



# intro2rev

You had to be there for the attendance flag :)

~RISC 27/8/25

# Acknowledgment of Country

RISC acknowledges the people of the Woi Wurrung and Boon Wurrung language groups of the eastern Kulin Nation on whose unceded lands we conduct the business of the University and the club. RISC acknowledges their Ancestors and Elders, past, present, and emerging





Today's Sponsor...

# BINARY NINJA



Binary Ninja is an interactive decompiler, disassembler, debugger, and binary analysis platform built by reverse engineers, for reverse engineers.

Developed with a focus on delivering a high-quality API for automation and a clean and usable GUI, Binary Ninja is in active use by malware analysts, vulnerability researchers, and software developers worldwide.

Decompile software built for many common architectures on Windows, macOS, and Linux.

# Hello, world!

## baby steps



- We've all (hopefully) written hello world before

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- Maybe in C?



```
#include <stdio.h>

int main(int argc, const char** argv) {
    printf("Hello, world!");
}
```

8,0-1 All

# Hello, world!

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- “Compile” using gcc and then run



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```

8,0-1 All



```
$ gcc hello_world.c -o hello_world
$ ./hello_world
Hello, world!
$
```

# Hello, world!

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- We've all (hopefully) written hello world before
- Maybe in C?
- “Compile” using gcc and then run
  - What actually happens here?



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- “Compile” using gcc and then run
  - What actually happens here?
  - Why can't I just run `hello_world.c`?



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8,0-1 All



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- “Compile” using gcc and then run
  - What actually happens here?
  - Why can't I just run `hello_world.c`?
  - What even is the `hello_world` file?



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```

8,0-1 All



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$ gcc hello_world.c -o hello_world
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Hello, world!
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```

# Hello, world!

## baby steps



- `file` command tells us a lot

```
$ file hello_world
hello_world: ELF 64-bit LSB pie executable, x86-64, version 1 (SYSV), dy
namically linked, interpreter /lib64/ld-linux-x86-64.so.2, BuildID[sha1]
=6fccaccbe66638ac049145ed4b8e591d443d5a4e, for GNU/Linux 3.2.0, not stri
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```

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  - Dynamically linked

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```

CPU...  
is

- *Compiled for x86-64 CPUs*



# CPU...

is



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- CPU → **C**entral **P**rocessing **U**nit

# CPU...

is



- *Compiled for x86-64 CPUs*
- CPU → **Central Processing Unit**
- Runs programs by executing **instructions**

# CPU...

is



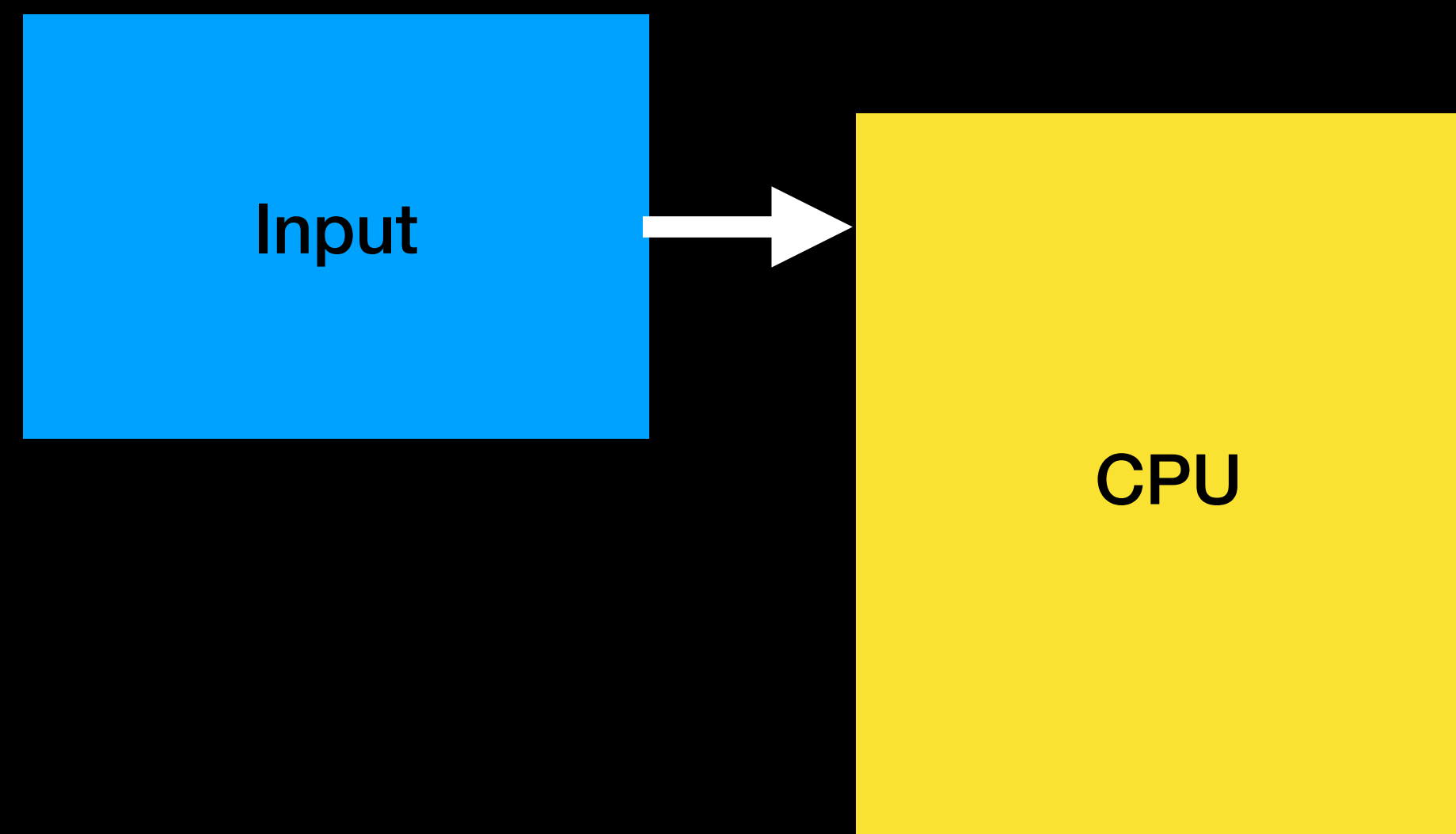
- *Compiled for x86-64 CPUs*
- CPU → **C**entral **P**rocessing **U**nit
- Runs programs by executing **instructions**
- One instruction at a time, **billions** per second (GHz)

# CPU...

as a machine



- CPU takes input (data)

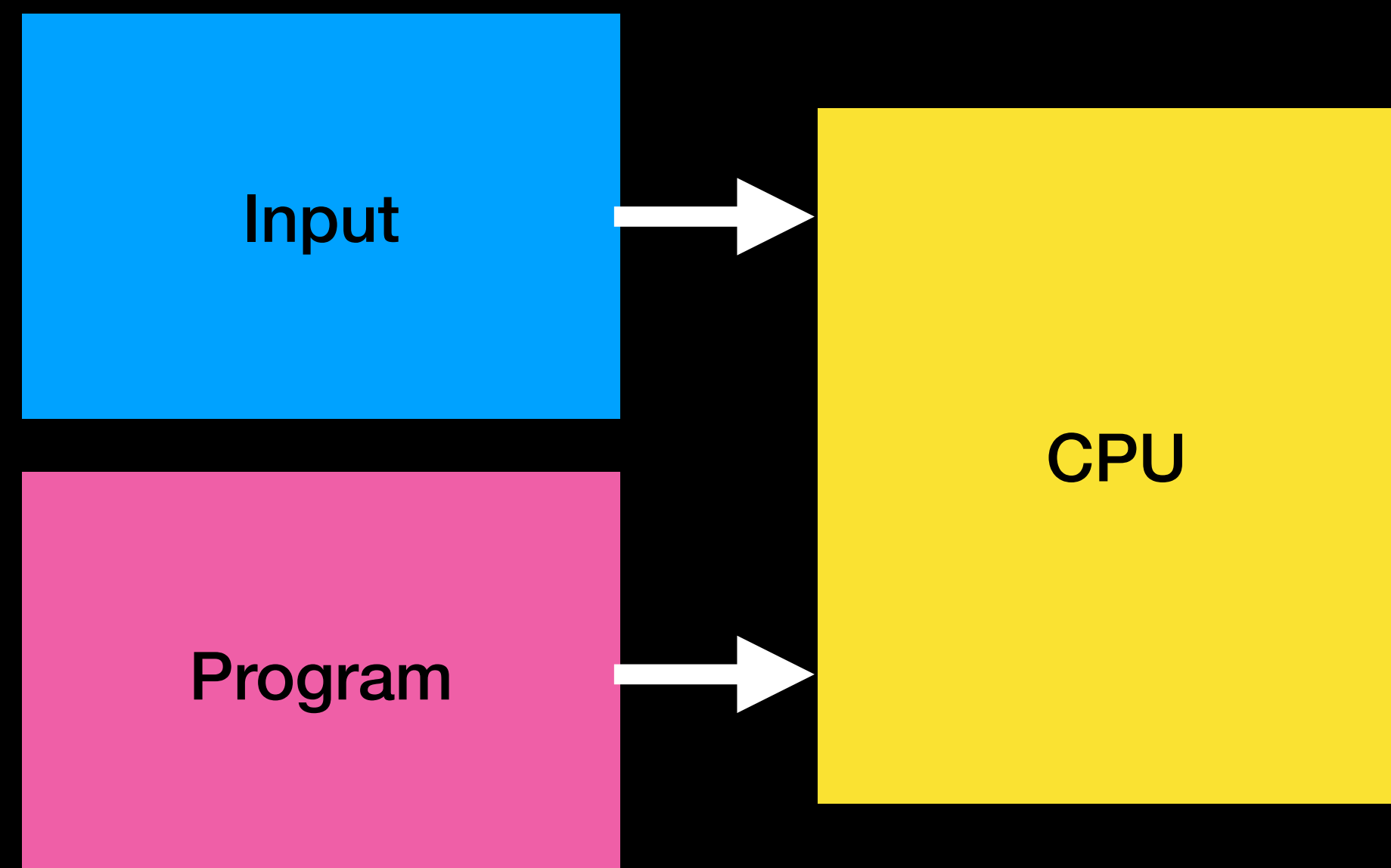


# CPU...

as a machine



- CPU takes **input (data)**
- Follows a **program (instructions)**

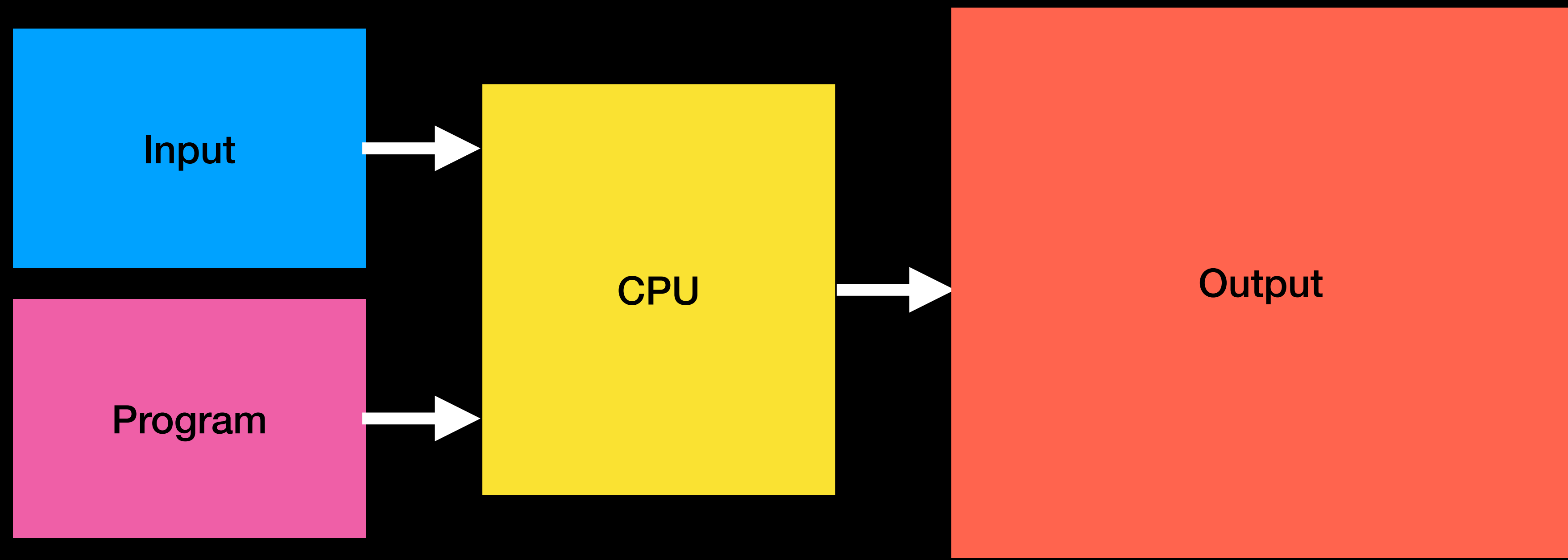


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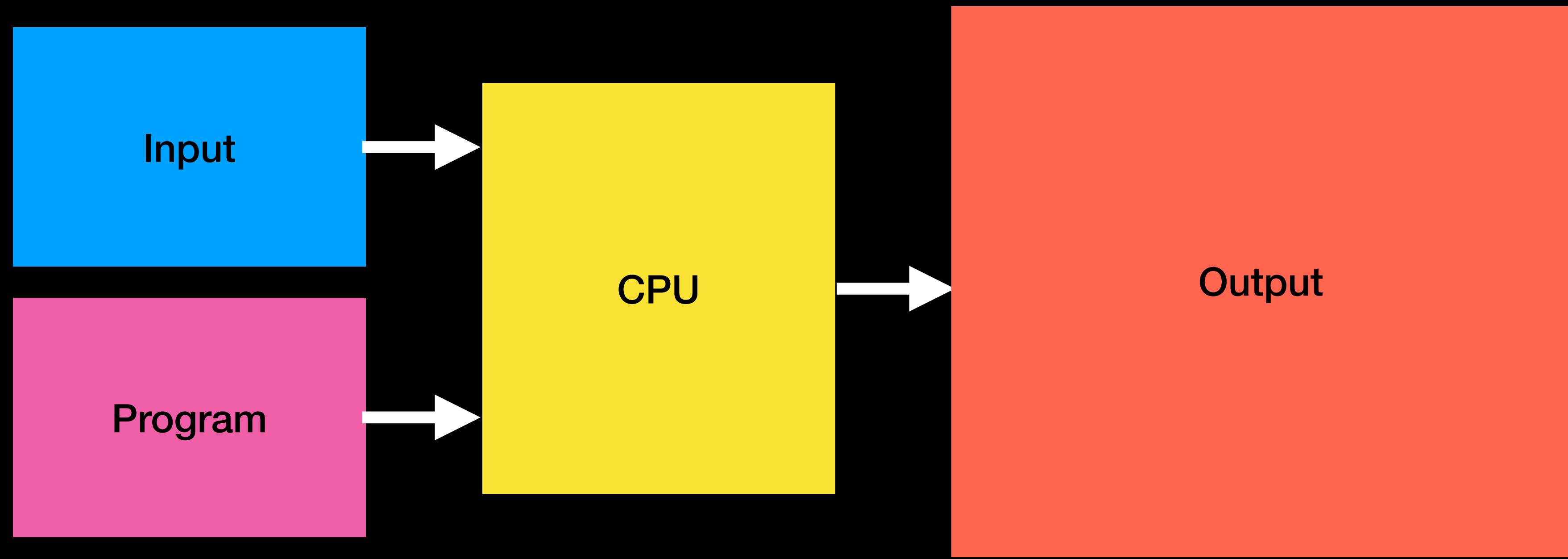


# CPU...

as a machine



- CPU takes **input (data)**
- Follows a **program (instructions)**
- Produces **output (results)**
- That's all!

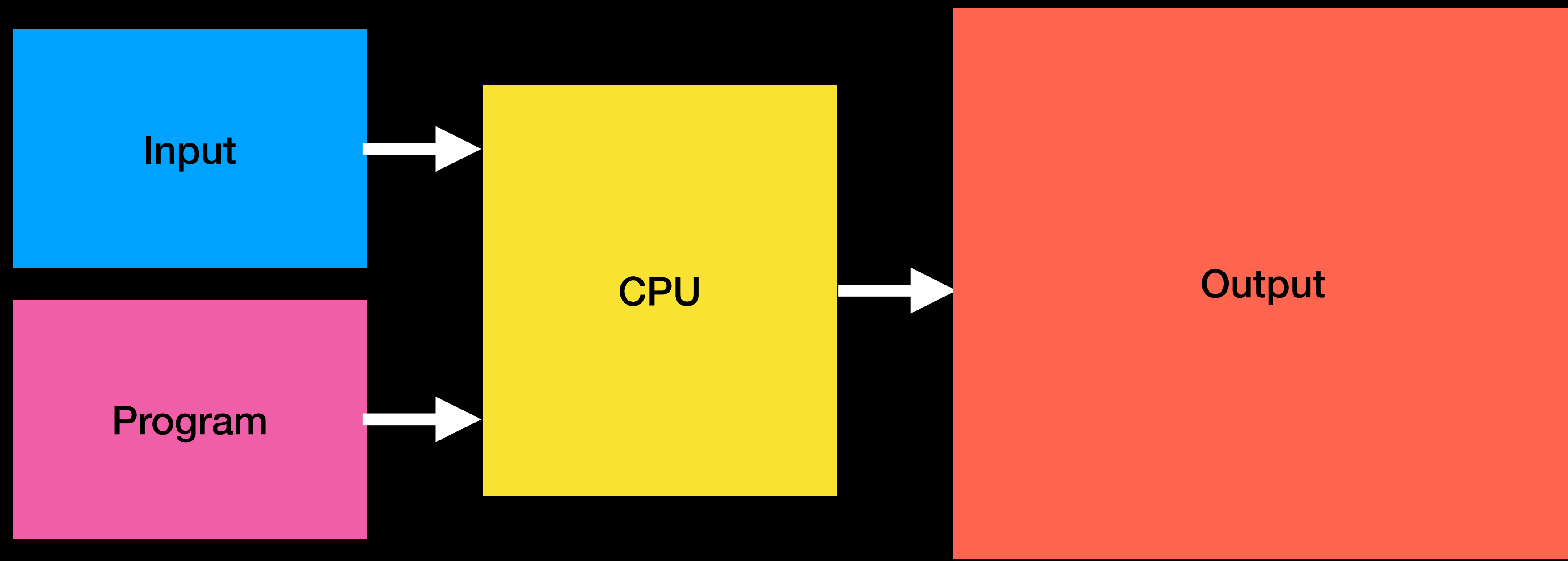


# CPU...

in reality



- Incredibly complex internal machinery

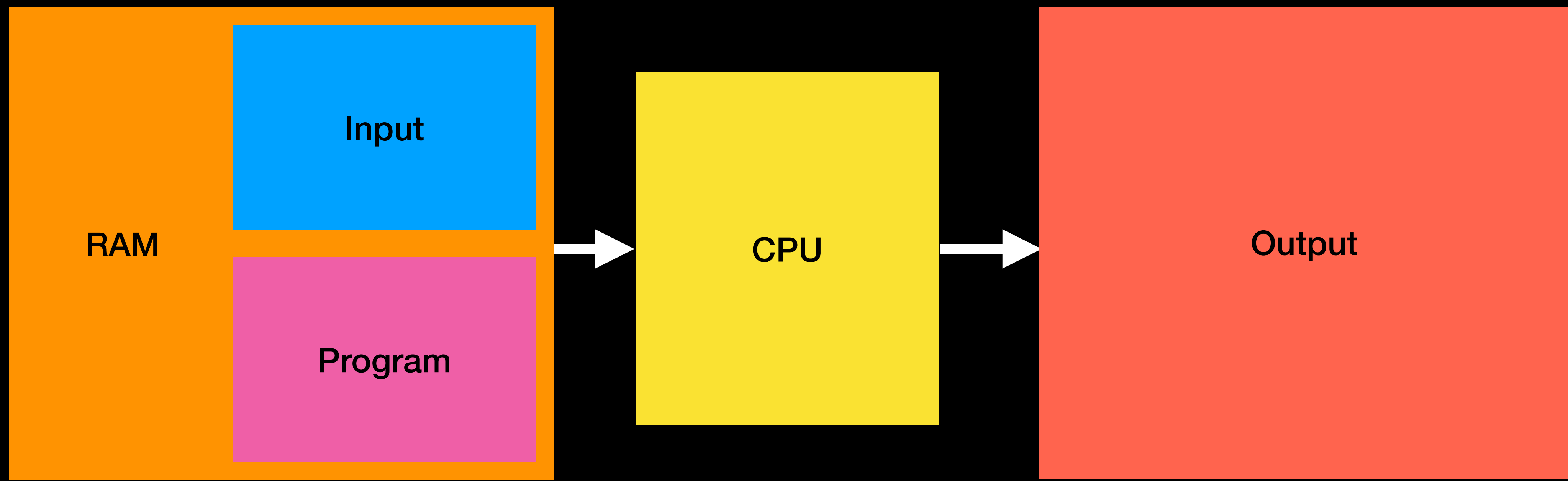


# CPU...

in reality



- Incredibly complex internal machinery
- Typically program and input refer to the same physical hardware (RAM)

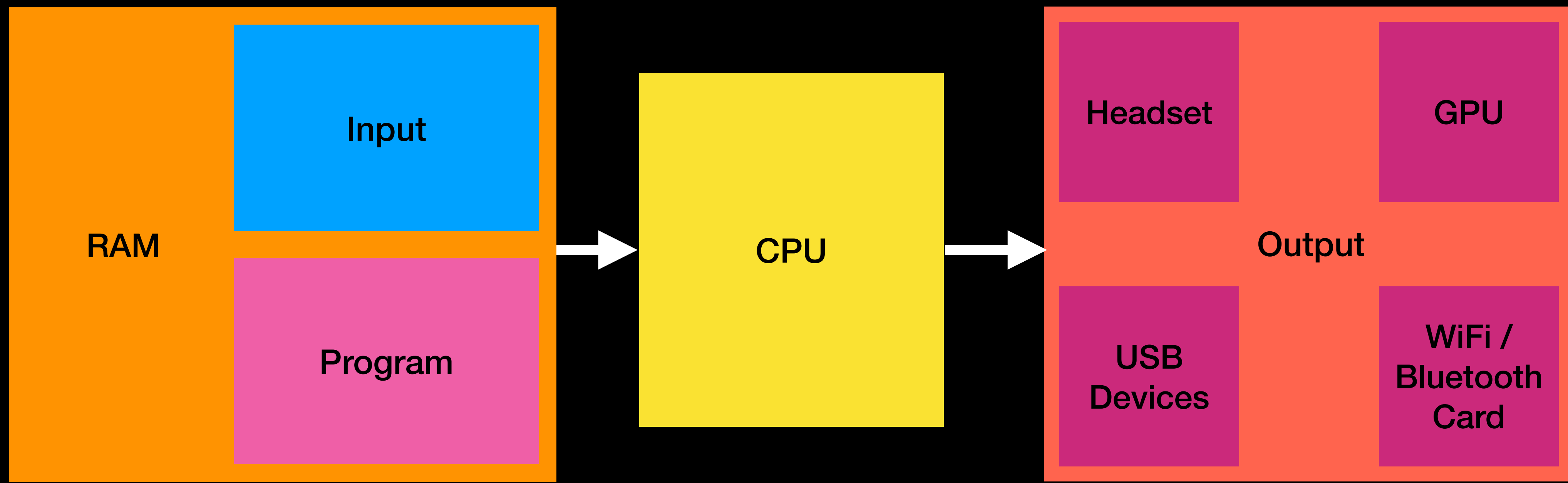


# CPU...

in reality



- Connect to multiple output devices (graphics, I/O devices, peripherals...)

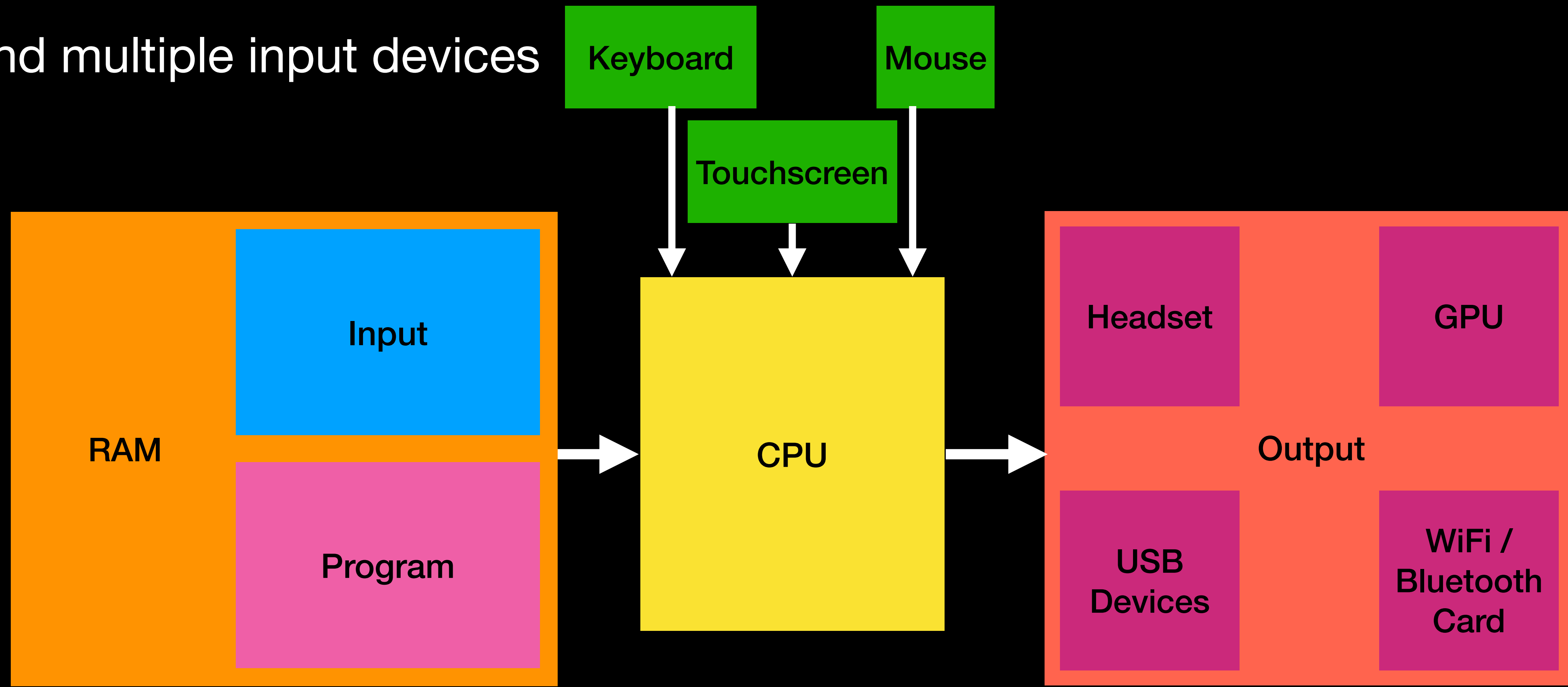


# CPU...

in reality



- Connect to multiple output devices (graphics, I/O devices, peripherals...)
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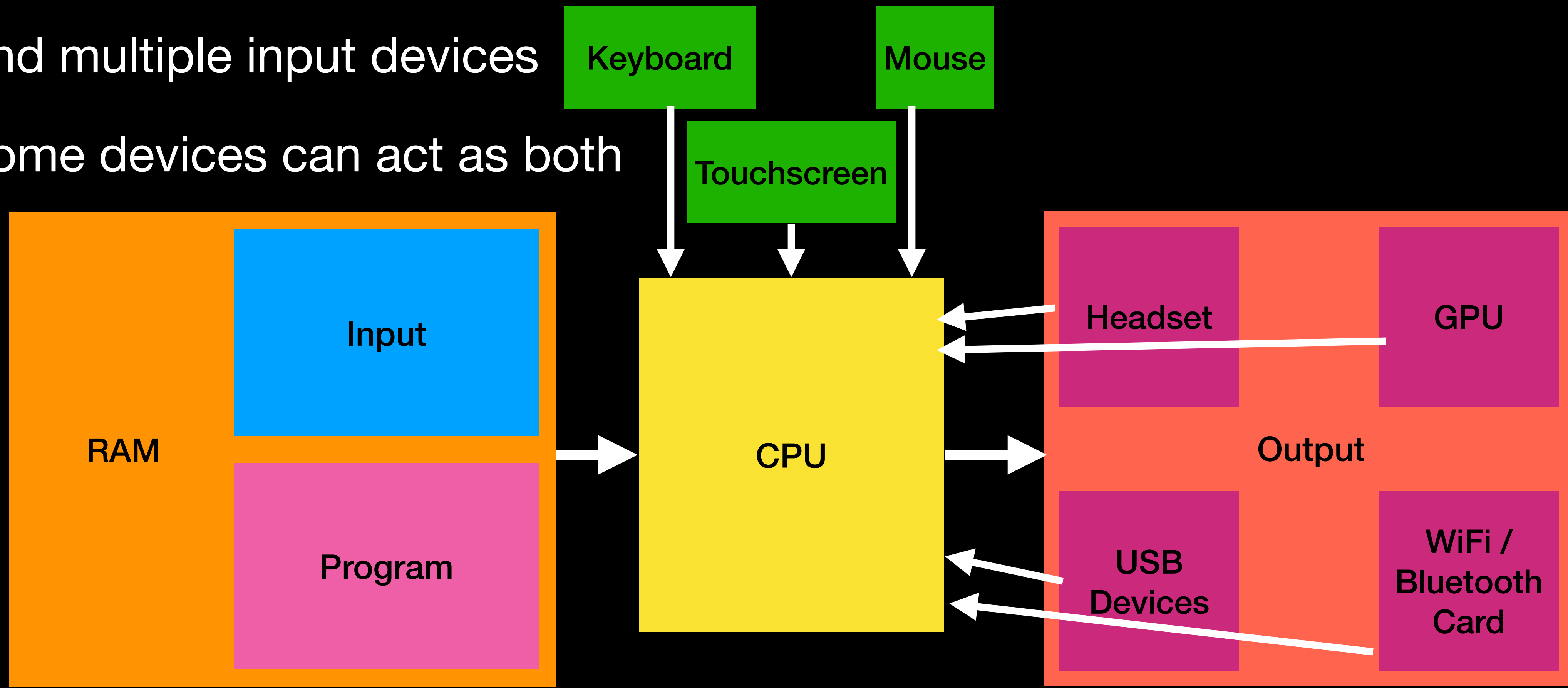


# CPU...

in reality



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- Some devices can act as both



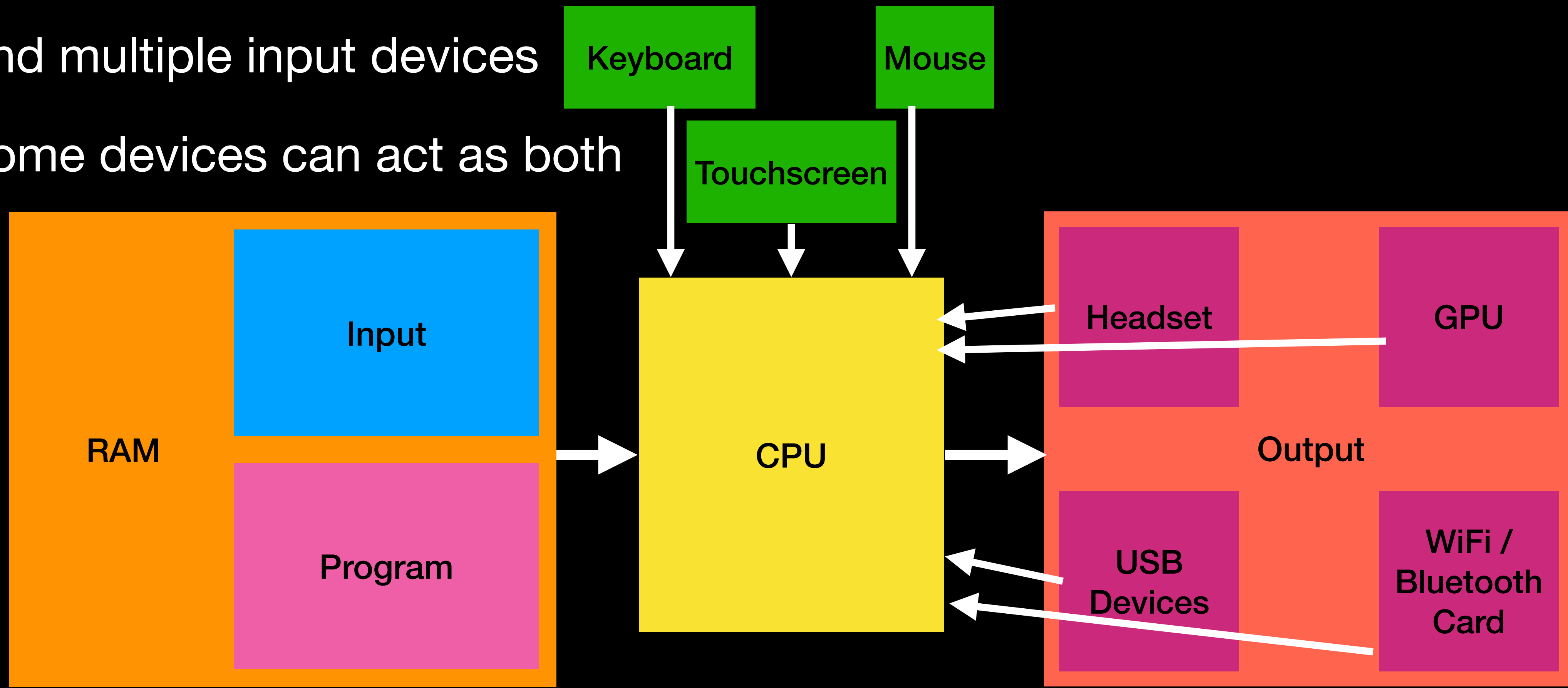
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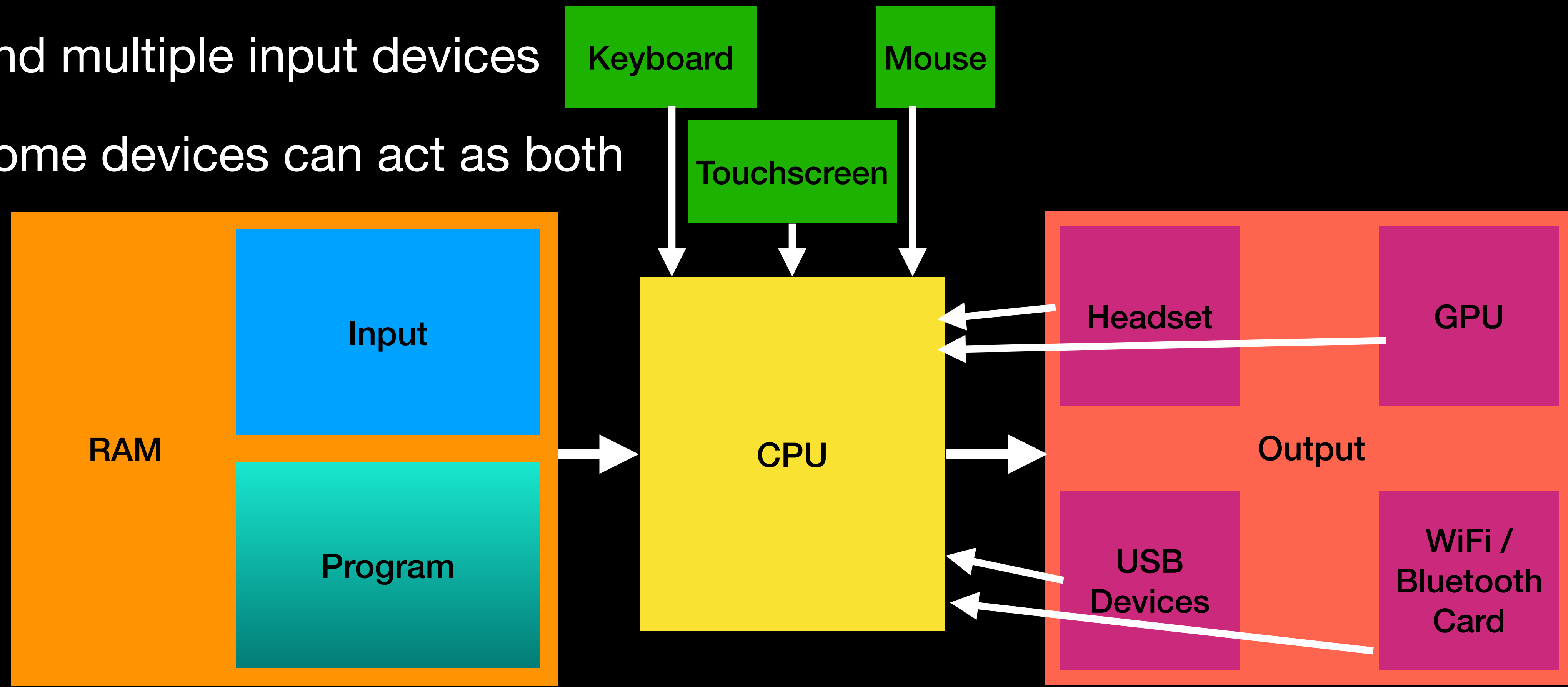
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- Different CPUs have different instruction “sets” (ISA)
  - We will focus on x86/x64 in this talk, but challenges may involve MIPS, ARM, ...

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# CPU...

has “registers”



- CPU keeps a set of internal “registers” that hold N-bit values

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- Some registers have special purposes (RIP, RSP, CS/DS/ES/..., EFLAGS, ...)
- Registers are much faster than RAM
  - Say we want to compute  $(x + 7) * 2 + 14$
  - Much faster to load X to a register and operate on that register

# CPU...

## instruction syntax

- Two main conventions:



# CPU...

## instruction syntax



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    - Also prefix registers with `%`, constants with `$`

# CPU...

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### Intel

```
mov  rax, 2
add  rax, 3
imul rax, 5
```

# CPU...

## instruction syntax



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  - AT&T Syntax - `<op> <src> <dst>`
  - Also prefix registers with `%`, constants with `$`

**Intel**

```
mov  rax, 2
add  rax, 3
imul rax, 5
```

**AT&T**

```
movq  $2, %rax
addq  $3, %rax
imulq $5, %rax
```

# CPU...

## instruction syntax



- Two main conventions:

- Intel Syntax - `<op> <dst> <src>`
- AT&T Syntax - `<op> <src> <dst>`

We will use Intel syntax for our workshops, but it is entirely personal preference

- Also prefix registers with `%`, constants with `$`

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mov  rax, 2
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### AT&T

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movq  $2, %rax
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# CPU...

## instruction syntax



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**Intel**  
`mov rax, 2`  
`add rax, 3`  
`imul rax, 5`

**What is the value of rax  
after this program?**

**AT&T**  
`movq $2, %rax`  
`addq $3, %rax`  
`imulq $5, %rax`

# CPU...

## instruction syntax



```
mov    rax, 2  
add    rax, 3  
imul   rax, 5
```

# CPU...

## instruction syntax



Move the value 2 into `rax`

```
mov    rax, 2  
add    rax, 3  
imul   rax, 5
```

`rax = 2`

# CPU...

## instruction syntax



Move the value 2 into `rax`

Add the value 3 to `rax`

```
mov    rax, 2  
add    rax, 3  
imul   rax, 5
```

`rax` = 2

`rax` = 5

# CPU...

## instruction syntax



Move the value 2 into `rax`

Add the value 3 to `rax`

Multiply `rax` by 5

```
mov    rax, 2
```

```
add    rax, 3
```

```
imul   rax, 5
```

`rax` = 2

`rax` = 5

`rax` = 25

# CPU...

## control flow



- Sometimes we want to do **X** if **Y**

# CPU...

## control flow



- Sometimes we want to do **X** if **Y**
  - **Make cake** if **have ingredients**

# CPU...

## control flow



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```
int foo(int x) {  
    if (x == 5) {  
        return 1;  
    } else {  
        return 0;  
    }  
}
```

-- INSERT -- 8,1 All

# CPU...

## control flow



- Sometimes we want to do **X** if **Y**
  - **Make cake** if **have ingredients**

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    if (x == 5) {  
        return 1;  
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}
```

-- INSERT -- 8,1 All



```
$ cc -c cfg.c -o cfg.o  
$ objdump -d -M intel cfg.o  
  
cfg.o:      file format elf64-x86-64  
  
Disassembly of section .text:  
  
0000000000000000 <foo>:  
0:  f3 0f 1e fa      endbr64  
4:  55               push    rbp  
5:  48 89 e5         mov     rbp, rsp  
8:  89 7d fc         mov     DWORD PTR [rbp-0x4], edi  
b:  83 7d fc 05      cmp     DWORD PTR [rbp-0x4], 0x5  
f:  75 07           jne     18 <foo+0x18>  
11: b8 01 00 00 00   mov     eax, 0x1  
16: eb 05           jmp     1d <foo+0x1d>  
18: b8 00 00 00 00   mov     eax, 0x0  
1d: 5d              pop     rbp  
1e: c3              ret  
$
```

# CPU...

## control flow



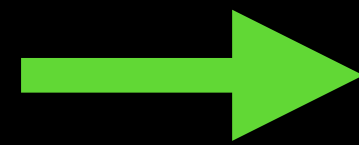
```
0:  endbr64
4:  push    rbp
5:  mov     rbp, rsp
8:  mov     DWORD PTR [rbp-0x4], edi
b:  cmp     DWORD PTR [rbp-0x4], 0x5
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1d: pop     rbp
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```

# CPU...

## control flow



Function entry



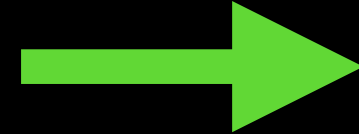
```
0: endbr64
4: push    rbp
5: mov     rbp, rsp
8: mov     DWORD PTR [rbp-0x4], edi
b: cmp     DWORD PTR [rbp-0x4], 0x5
f: jne     18 <foo+0x18>
11: mov     eax, 0x1
16: jmp     1d <foo+0x1d>
18: mov     eax, 0x0
1d: pop     rbp
1e: ret
```

# CPU...

## control flow



Set up stack  
(more on this later!)



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```

# CPU...

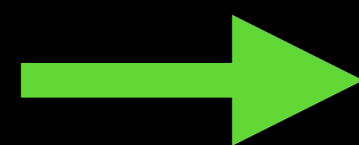
## control flow



```
int foo(int x) {  
    if (x == 5) {  
        return 1;  
    } else {  
        return 0;  
    }  
}
```

-- INSERT -- 8,1 A11

Compare x to 5



```
0:  endbr64  
4:  push    rbp  
5:  mov     rbp, rsp  
8:  mov     DWORD PTR [rbp-0x4], edi  
b:  cmp     DWORD PTR [rbp-0x4], 0x5  
f:  jne     18 <foo+0x18>  
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1d: pop     rbp  
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# CPU...

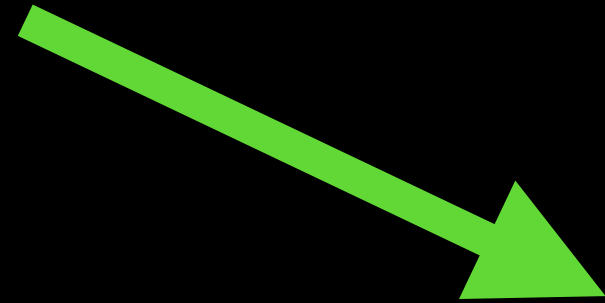
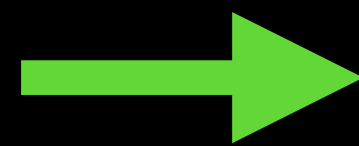
## control flow



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}
```

-- INSERT -- 8,1 A11

If not equal (ne), jump (j)  
to instruction at 0x18



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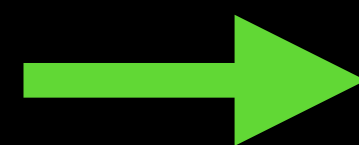
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    }  
}
```

-- INSERT -- 8,1 A11

True branch



```
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11: mov     eax, 0x1  
16: jmp     1d <foo+0x1d>  
18: mov     eax, 0x0  
1d: pop     rbp  
1e: ret
```

# CPU...

## control flow



```
int foo(int x) {  
    if (x == 5) {  
        return 1;  
    } else {  
        return 0;  
    }  
}
```

-- INSERT -- 8,1 A11

Jump to instruction  
at address 0x1d

```
0:  endbr64  
4:  push    rbp  
5:  mov     rbp, rsp  
8:  mov     DWORD PTR [rbp-0x4], edi  
b:  cmp     DWORD PTR [rbp-0x4], 0x5  
f:  jne     18 <foo+0x18>  
11: mov     eax, 0x1  
16: jmp     1d <foo+0x1d>  
18: mov     eax, 0x0  
1d: pop     rbp  
1e: ret
```

# CPU...

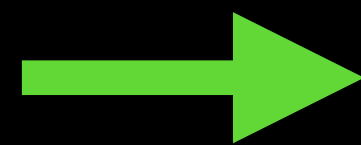
## control flow



```
int foo(int x) {  
    if (x == 5) {  
        return 1;  
    } else {  
        return 0;  
    }  
}
```

-- INSERT -- 8,1 A11

False branch



```
0:  endbr64  
4:  push    rbp  
5:  mov     rbp, rsp  
8:  mov     DWORD PTR [rbp-0x4], edi  
b:  cmp     DWORD PTR [rbp-0x4], 0x5  
f:  jne     18 <foo+0x18>  
11: mov     eax, 0x1  
16: jmp     1d <foo+0x1d>  
18: mov     eax, 0x0  
1d: pop     rbp  
1e: ret
```

# CPU...

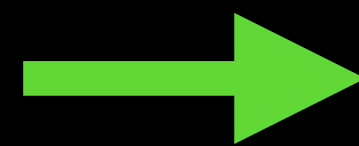
## control flow



```
int foo(int x) {  
    if (x == 5) {  
        return 1;  
    } else {  
        return 0;  
    }  
}
```

-- INSERT -- 8,1 A11

Leave function  
(return value is  
stored in rax)



```
0: endbr64  
4: push    rbp  
5: mov     rbp, rsp  
8: mov     DWORD PTR [rbp-0x4], edi  
b: cmp     DWORD PTR [rbp-0x4], 0x5  
f: jne     18 <foo+0x18>  
11: mov     eax, 0x1  
16: jmp     1d <foo+0x1d>  
18: mov     eax, 0x0  
1d: pop     rbp  
1e: ret
```

# CPU...

## control flow

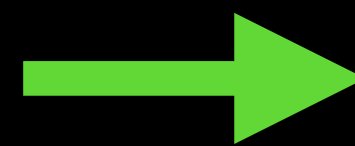


```
int foo(int x) {  
    if (x == 5) {  
        return 1;  
    } else {  
        return 0;  
    }  
}
```

-- INSERT -- 8,1 A11

**This kinda sucks to read, doesn't it?**

Leave function  
(return value is  
stored in `rax`)



```
0:  endbr64  
4:  push    rbp  
5:  mov     rbp, rsp  
8:  mov     DWORD PTR [rbp-0x4], edi  
b:  cmp     DWORD PTR [rbp-0x4], 0x5  
f:  jne     18 <foo+0x18>  
11: mov     eax, 0x1  
16: jmp     1d <foo+0x1d>  
18: mov     eax, 0x0  
1d: pop     rbp  
1e: ret
```

# CPU...

## control flow graphs



```
0:  endbr64
4:  push    rbp
5:  mov     rbp, rsp
8:  mov     DWORD PTR [rbp-0x4], edi
b:  cmp     DWORD PTR [rbp-0x4], 0x5
f:  jne     18 <foo+0x18>
```

```
11: mov     eax, 0x1
16: jmp     1d <foo+0x1d>
```

```
18: mov     eax, 0x0
```

```
1d: pop     rbp
1e: ret
```

# CPU...

## control flow graphs



**Graph view lets  
us understand  
the program in  
small isolated  
sections**

```
0:  endbr64
4:  push    rbp
5:  mov     rbp, rsp
8:  mov     DWORD PTR [rbp-0x4], edi
b:  cmp     DWORD PTR [rbp-0x4], 0x5
f:  jne     18 <foo+0x18>
```

```
11: mov     eax, 0x1
16: jmp     1d <foo+0x1d>
```

```
18: mov     eax, 0x0
```

```
1d: pop     rbp
1e: ret
```

# CPU...

## control flow graphs



**Graph view lets  
us understand  
the program in  
small isolated  
sections**

```
0:  endbr64
4:  push    rbp
5:  mov     rbp, rsp
8:  mov     DWORD PTR [rbp-0x4], edi
b:  cmp     DWORD PTR [rbp-0x4], 0x5
f:  jne     18 <foo+0x18>
```

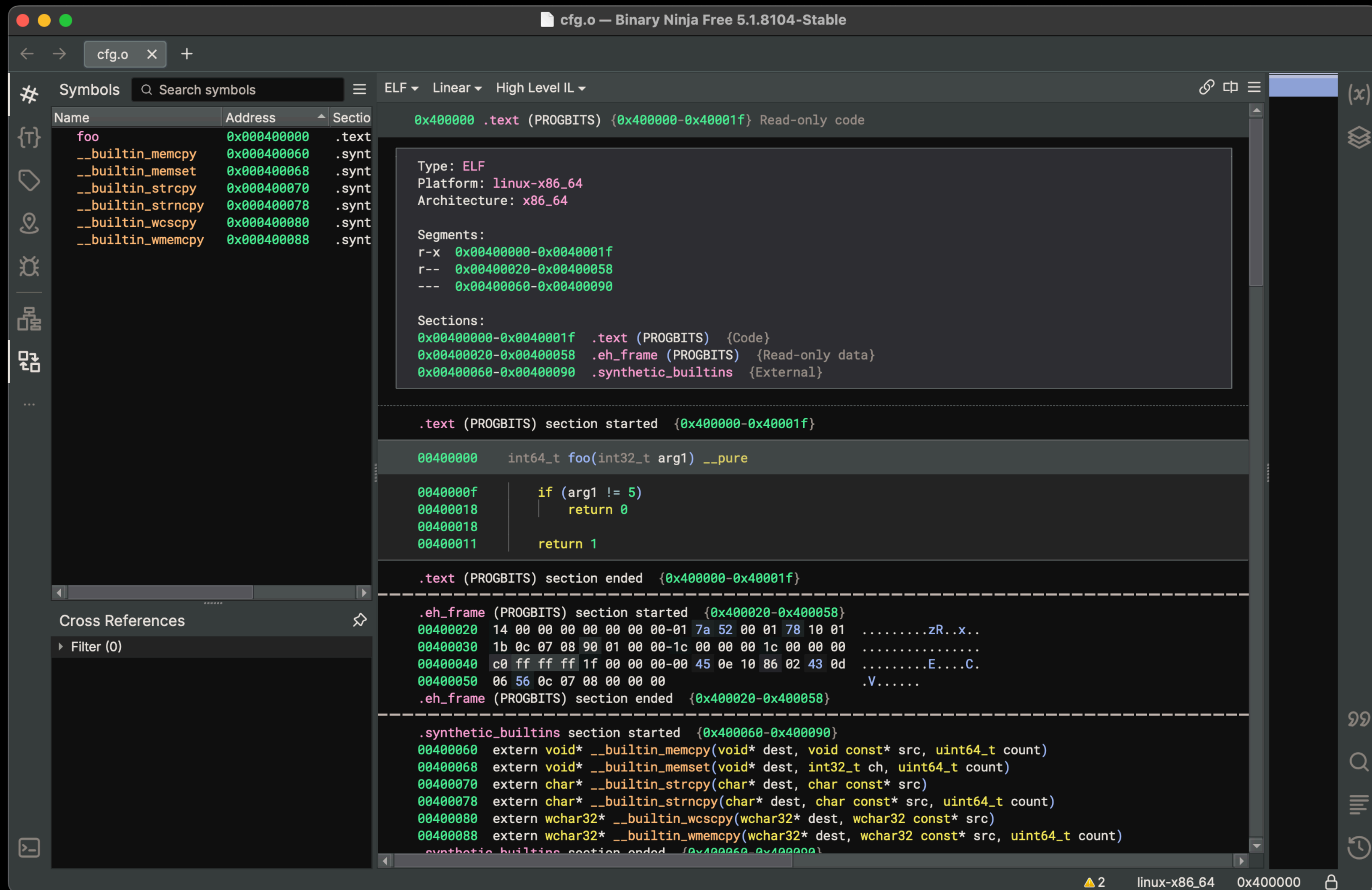
**Binary Ninja can  
do this for us :)**

```
11: mov     eax, 0x1
16: jmp     1d <foo+0x1d>
```

```
18: mov     eax, 0x0
```

```
1d: pop     rbp
1e: ret
```

# Binary Ninja



# Binary Ninja



cfgo

cfgo

+

# Symbols

Search symbols

...

Name	Address	Section
foo	0x00040000	.text
__builtin_memcpy	0x00040060	.synt
__builtin_memset	0x00040068	.synt
__builtin_strcpy	0x00040070	.synt
__builtin_strncpy	0x00040078	.synt
__builtin_wscpy	0x00040080	.synt
__builtin_wmemcpy	0x00040088	.synt

Cross References

Filter (1)

Code References

foo

0040000f jne 0x400018

ELF

Linear

Disassembly

Disassembly

Low Level IL

Medium Level IL

High Level IL

Pseudo C

Pseudo Objective-C

Advanced IL Forms

int64\_t foo(int32\_t arg1) \_\_pure

00400000 f30f1efa endbr64

00400004 55 push rbp {\_\_saved\_rbp}

00400005 4889e5 mov rbp, rsp {\_\_saved\_rbp}

00400008 897dfc mov dword [rbp-0x4 {var\_c}], edi

0040000b 837dfc05 cmp dword [rbp-0x4 {var\_c}], 0x5

0040000f 7507 jne 0x400018

00400011 b801000000 mov eax, 0x1

00400016 eb05 jmp 0x40001d

00400018 b800000000 mov eax, 0x0

0040001d 5d pop rbp {\_\_saved\_rbp}

0040001e c3 retn {\_\_return\_addr}

.text (PROGBITS) section ended {0x400000-0x40001f}

.eh\_frame (PROGBITS) section started {0x400020-0x400058}

00400020 14 00 00 00 00 00 00-01 7a 52 00 01 78 10 01 .....zR..X..

00400030 1b 0c 07 08 90 01 00 00-1c 00 00 00 1c 00 00 00 .....

00400040 c0 ff ff ff 1f 00 00 00-00 45 0e 10 86 02 43 0d .....E....C.

00400050 06 56 0c 07 08 00 00 00 .....V.....

.eh\_frame (PROGBITS) section ended {0x400020-0x400058}

.synthetic\_builtins section started {0x400060-0x400090}

00400060 extern void\* \_\_builtin\_memcpy(void\* dest, void const\* src, uint64\_t count)

00400068 extern void\* \_\_builtin\_memset(void\* dest, int32\_t ch, uint64\_t count)

00400070 extern char\* \_\_builtin\_strcpy(char\* dest, char const\* src)

00400078 extern char\* \_\_builtin\_strncpy(char\* dest, char const\* src, uint64\_t count)

00400080 extern wchar\_t\* \_\_builtin\_wscpy(wchar\_t\* dest, wchar\_t const\* src)

2

linux-x86\_64

0x400018-0x40001d (0x5 bytes)

🔒

# Binary Ninja



cfg.o — Binary Ninja Free 5.1.8104-Stable

cfgo

Search symbols

Name	Address	Section
foo	0x00040000	.text
__builtin_memcpy	0x00040060	.synt
__builtin_memset	0x00040068	.synt
__builtin_strcpy	0x00040070	.synt
__builtin_strncpy	0x00040078	.synt
__builtin_wscpy	0x00040080	.synt
__builtin_wmemcpy	0x00040088	.synt

ELF

Graph

Byte Overview

Graph

Hex Editor

Linear

Memory Map

Strings

Triage Summary

Types

Workflow

foo:

00400000 endbr64

00400004 push rbp {\_\_saved\_rbp}

00400005 mov rbp, rsp {\_\_saved\_rbp}

00400008 mov dword [rbp-0x4 {var\_c}], edi

0040000b cmp dword [rbp-0x4 {var\_c}], 0x5

0040000f jne 0x400018

00400018 mov eax, 0x0

00400011 mov eax, 0x1

00400016 jmp 0x40001d

0040001d pop rbp {\_\_saved\_rbp}

0040001e retn {\_\_return\_addr}

Cross References

Filter (3)

Code References {2}

foo {2}

00400011 mov eax, 0x1

00400018 mov eax, 0x0

Variable References {1}

int32\_t var\_c {1}

0040000f jne 0x400018

2

linux-x86\_64

0x40000f-0x400011 (0x2 bytes)



What if we could turn  
assembly → C style code?

# Binary Ninja



cfg.o — Binary Ninja Free 5.1.8104-Stable

cfg.o x +

# Symbols

Search symbols

Name	Address	Section
foo	0x00040000	.text
__builtin_memcpy	0x00040060	.synt
__builtin_memset	0x00040068	.synt
__builtin_strcpy	0x00040070	.synt
__builtin_strncpy	0x00040078	.synt
__builtin_wscpy	0x00040080	.synt
__builtin_wmemcpy	0x00040088	.synt

ELF Linear Pseudo C

0x400000 .text (PROGBITS) {0x400000-0x40001f} Read-only code

0x00400000-0x0040001f .text (PROGBITS) {Code}

0x00400020-0x00400058 .eh\_frame (PROGBITS) {Read-only data}

0x00400060-0x00400090 .synthetic\_builtins {External}

.text (PROGBITS) section started {0x400000-0x40001f}

00400000 int64\_t foo(int32\_t arg1) \_\_pure

00400000 {

00400000 if (arg1 != 5)

00400018 | return 0;

00400018

00400018 return 1;

00400011

00400000 }

.text (PROGBITS) section ended {0x400000-0x40001f}

.eh\_frame (PROGBITS) section started {0x400020-0x400058}

00400020 14 00 00 00 00 00 00 00-01 7a 52 00 01 78 10 01 .....zR..x..

00400030 1b 0c 07 08 90 01 00 00-1c 00 00 00 1c 00 00 00 .....

00400040 c0 ff ff ff 1f 00 00 00-00 45 0e 10 86 02 43 0d .....E....C.

00400050 06 56 0c 07 08 00 00 00 .....V.....

00400050 .eh\_frame (PROGBITS) section ended {0x400020-0x400058}

.synthetic\_builtins section started {0x400060-0x400090}

00400060 extern void\* \_\_builtin\_memcpy(void\* dest, void const\* src, uint64\_t count)

00400068 extern void\* \_\_builtin\_memset(void\* dest, int32\_t ch, uint64\_t count)

00400070 extern char\* \_\_builtin\_strcpy(char\* dest, char const\* src)

00400078 extern char\* \_\_builtin\_strncpy(char\* dest, char const\* src, uint64\_t count)

00400080 extern wchar32\* \_\_builtin\_wscpy(wchar32\* dest, wchar32 const\* src)

00400088 extern wchar32\* \_\_builtin\_wmemcpy(wchar32\* dest, wchar32 const\* src, uint64\_t count)

.synthetic\_builtins section ended {0x400060-0x400090}

Cross References

Filter (0)

2 linux-x86\_64 0x400000-0x400004 (0x4 bytes)

# Binary Ninja



cfg.o — Binary Ninja Free 5.1.8104-Stable

cfgo x +

# Symbols Search symbols

ELF Linear Pseudo C

Name	Address	Section
foo	0x00040000	.text
__builtin_memcpy	0x00040060	.synt
__builtin_memset	0x00040068	.synt
__builtin_strcpy	0x00040070	.synt
__builtin_strncpy	0x00040078	.synt
__builtin_wscpy	0x00040080	.synt
__builtin_wmemcpy	0x00040088	.synt

0x400000 .text (PROGBITS) {0x400000-0x40001f} Read-only code

0x00400000-0x0040001f .text (PROGBITS) {0x00400000-0x0040001f} Read-only code

0x00400020-0x00400058 .eh\_frame (PROGBITS) {0x00400020-0x00400058} Read-only code

0x00400060-0x00400090 .synthetic\_builtins {0x00400060-0x00400090} External code

.text (PROGBITS) section started {0x400000-0x40001f}

```
00400000 int64_t foo(int32_t arg1) __pure
00400000 {
00400000     if (arg1 != 5)
00400018         return 0;
00400018
00400018     return 1;
00400000 }
```

.text (PROGBITS) section ended {0x400000-0x40001f}

.eh\_frame (PROGBITS) section started {0x400020-0x400058}

```
00400020 14 00 00 00 00 00 00 01 7a 52 00 01 78
00400030 1b 0c 07 08 90 01 00 00 1c 00 00 00 1c 00
00400040 c0 ff ff ff 1f 00 00 00 45 0e 10 86 02
00400050 06 56 0c 07 08 00 00 00
.eh_frame (PROGBITS) section ended {0x400020-0x400058}

.synthetic_builtins section started {0x400060-0x400090}



```
00400060 extern void* __builtin_memcpy(void* dest, void const* src, uint64_t count)
00400068 extern void* __builtin_memset(void* dest, int32_t ch, uint64_t count)
00400070 extern char* __builtin_strcpy(char* dest, char const* src)
00400078 extern char* __builtin_strncpy(char* dest, char const* src, uint64_t count)
00400080 extern wchar32* __builtin_wscpy(wchar32* dest, wchar32 const* src)
00400088 extern wchar32* __builtin_wmemcpy(wchar32* dest, wchar32 const* src, uint64_t count)
.synthetic_builtins section ended {0x400060-0x400090}

Cross References



Filter (0)



```
int64_t foo(int32_t arg1) __pure
{
    if (arg1 != 5)
        return 0;

    return 1;
}
```



2 linux-x86_64 0x400000-0x400004 (0x4 bytes)


```


```

# Binary Ninja

## Decompilation



- Key things to note:

```
int64_t foo(int32_t arg1) __pure
{
    if (arg1 != 5)
        return 0;

    return 1;
}
```

# Binary Ninja

## Decompilation



- Key things to note:
  - Variable names are lost

```
int64_t foo(int32_t arg1) __pure
{
    if (arg1 != 5)
        return 0;

    return 1;
}
```

# Binary Ninja

## Decompilation



- Key things to note:
  - Variable names are lost
  - Exact source structure is lost

```
int64_t foo(int32_t arg1) __pure
{
    if (arg1 != 5)
        return 0;

    return 1;
}
```

# Binary Ninja

## Decompilation



- Key things to note:
  - Variable names are lost
  - Exact source structure is lost
  - Comments are lost

```
int64_t foo(int32_t arg1) __pure
{
    if (arg1 != 5)
        return 0;

    return 1;
}
```

# Binary Ninja

## Decompilation



- Key things to note:
  - Variable names are lost
  - Exact source structure is lost
  - Comments are lost
  - Function names *may* be lost

```
int64_t foo(int32_t arg1) __pure
{
    if (arg1 != 5)
        return 0;

    return 1;
}
```

# Binary Ninja

## Decompilation



- Key things to note:
  - Variable names are lost
  - Exact source structure is lost
  - Comments are lost
  - Function names *may* be lost
  - ***Decompilation is not an exact science***

```
int64_t foo(int32_t arg1) __pure
{
    if (arg1 != 5)
        return 0;

    return 1;
}
```

# Binary Ninja

## Decompilation



- Key things to note:
  - Variable names are lost
  - Exact source structure is lost
  - Comments are lost
  - Function names *may* be lost
  - ***Decompilation is not an exact science***
  - But it is still immensely helpful to understand how a program works :)

```
int64_t foo(int32_t arg1) __pure
{
    if (arg1 != 5)
        return 0;

    return 1;
}
```



**What if function names are  
missing?**

# Reverse Engineering

Making sense of things



- I've “stripped” a binary of it's symbols, removing any function name info

# Rev Makin

- I've “

pw\_s — Binary Ninja Free 5.1.8104-Stable

pw\_s x +

Symbols Search symbols

Name	Address	Sec
sub_401080	0x000401080	.pl
sub_401090	0x000401090	.pl
sub_4010a0	0x0004010a0	.pl
sub_4010b0	0x0004010b0	.pl
__cxa_finalize	0x0004010c0	.pl
strncpy	0x0004010d0	.pl
puts	0x0004010e0	.pl
strlen	0x0004010f0	.pl
__stack_chk_fail	0x000401100	.pl
printf	0x000401110	.pl
strcspn	0x000401120	.pl
fgets	0x000401130	.pl
strcmp	0x000401140	.pl
__ctype_b_loc	0x000401150	.pl
_start	0x000401160	.te
deregister_tm_clon...	0x000401190	.te
sub_4011c0	0x0004011c0	.te
_FINI_0	0x000401200	.te
_INIT_0	0x000401240	.te
sub_401249	0x000401249	.te
sub_4012db	0x0004012db	.te
sub_401364	0x000401364	.te
main	0x0004013eb	.te
_fini	0x0004014a4	.fi
strncpy	0x000403f90	.go
puts	0x000403f98	.go
strlen	0x000403fa0	.go

Cross References

Filter (1)

Code References {1}

main {1}

← 00401460 if (!sub\_401364(&buf))

ELF Linear Pseudo C

```
int32_t main(int32_t argc, char** argv, char** envp)
004013e4 /* no return */
00401364 }

004013eb int32_t main(int32_t argc, char** argv, char** envp)
{
    void* fsbase;
    int64_t rax = *(uint64_t*)((char*)fsbase + 0x28);
    printf("Enter password: ");
    char buf[0x48];
    int32_t result;

    if (fgets(&buf, 0x40, stdin))
    {
        buf[strcspn(&buf, "\n")] = 0;

        if (!sub_401364(&buf))
            puts("Wrong!");
        else
            puts("Correct!");

        result = 0;
    }
    else
        result = 1;

    *(uint64_t*)((char*)fsbase + 0x28);

    if (rax == *(uint64_t*)((char*)fsbase + 0x28))
        return result;

    __stack_chk_fail();
    /* no return */
}

.text (PROGBITS) section ended {0x401160-0x4014a4}

.fini (PROGBITS) section started {0x4014a4-0x4014b1}

004014a4 int64_t _fini() __pure
```

linux-x86\_64 0x401364-0x401368 (0x4 bytes)



```

0x0004010b0 .pl
0x0004010c0 .pl
0x0004010d0 .pl
0x0004010e0 .pl
0x0004010f0 .pl
fail 0x000401100 .pl
0x000401110 .pl
0x000401120 .pl
0x000401130 .pl
0x000401140 .pl
0x000401150 .pl
0x000401160 .te
m_clon... 0x000401190 .te
0x0004011c0 .te
0x000401200 .te
0x000401240 .te
0x000401249 .te
0x0004012db .te
0x000401364 .te
0x0004013eb .te
0x0004014a4 .fi
0x000403f90 .go
0x000403f98 .go
0x000403fa0 .no

```

```

es
es {1}
{1}
if (!sub_401364(&buf))

```

```

004013eb int32_t main(int32_t argc, char** argv, char** envp)

004013eb {
004013f7 void* fsbase;
00401415 int64_t rax = *(uint64_t*)((char*)fsbase + 0x28);
00401435 printf("Enter password: ");
00401435 char buf[0x48];
00401435 int32_t result;

00401435 if (fgets(&buf, 0x40, stdin))
00401435 {
00401454     buf[strcspn(&buf, "\n")] = 0;
00401454
00401467     if (!sub_401364(&buf))
00401484         puts("Wrong!");
00401467     else
00401473         puts("Correct!");
00401473
00401489     result = 0;
00401435 }
00401435 else
00401437     result = 1;
00401437

00401492 *(uint64_t*)((char*)fsbase + 0x28);
00401492

0040149b if (rax == *(uint64_t*)((char*)fsbase + 0x28))
004014a3     return result;
004014a3

0040149d __stack_chk_fail();
0040149d /* no return */
004013eb }

```

```

text (PROGBITS) section ended at 0x401160-0x4011a4

```

0x0004010b0 .pl  
0x0004010c0 .pl  
0x0004010d0 .pl  
0x0004010e0 .pl  
0x0004010f0 .pl  
fail 0x000401100 .pl  
0x000401110 .pl  
0x000401120 .pl  
0x000401130 .pl  
0x000401140 .pl  
0x000401150 .pl  
0x000401160 .te  
m\_clon... 0x000401190 .te  
0x0004011c0 .te  
0x000401200 .te  
0x000401240 .te  
0x000401249 .te  
0x0004012db .te  
0x000401364 .te  
0x0004013eb .te  
0x0004014a4 .fi  
0x000403f90 .go  
0x000403f98 .go  
0x000403fa0 .no

es {1}  
es {1}  
if (!sub\_401364(&buf))

```
004013eb int32_t main(int32_t argc, char** argv, char** envp)
{
    void* fsbase;
    int64_t rax = *(uint64_t*)((char*)fsbase + 0x28);
    printf("Enter password: ");
    char buf[0x48];
    int32_t result;

    if (fgets(&buf, 0x40, stdin))
    {
        buf[strcspn(&buf, "\n")] = 0;

        if (!sub_401364(&buf))
            puts("Wrong!");
        else
            puts("Correct!");

        result = 0;
    }
    else
        result = 1;

    *(uint64_t*)((char*)fsbase + 0x28);

    if (rax == *(uint64_t*)((char*)fsbase + 0x28))
        return result;

    __stack_chk_fail();
    /* no return */
}
```

Some functions are “imported”  
from other places

We will still have names for these

0x0004010b0 .pl  
0x0004010c0 .pl  
0x0004010d0 .pl  
0x0004010e0 .pl  
0x0004010f0 .pl  
fail 0x000401100 .pl  
0x000401110 .pl  
0x000401120 .pl  
0x000401130 .pl  
0x000401140 .pl  
0x000401150 .pl  
0x000401160 .te  
m\_clon... 0x000401190 .te  
0x0004011c0 .te  
0x000401200 .te  
0x000401240 .te  
0x000401249 .te  
0x0004012db .te  
0x000401364 .te  
0x0004013eb .te  
0x0004014a4 .fi  
0x000403f90 .go  
0x000403f98 .go  
0x000403fa0 .no

es {1}  
es {1}  
if (!sub\_401364(&buf))

```
004013eb  int32_t main(int32_t argc, char** argv, char** envp)
004013eb  {
004013f7      void* fsbase;
00401415      int64_t rax = *(uint64_t*)((char*)fsbase + 0x28);
00401435      printf("Enter password: ");
00401435      char buf[0x48];
00401435      int32_t result;
00401435
00401435      if (fgets(&buf, 0x40, stdin))
00401435      {
00401454          buf[strcspn(&buf, "\n")] = 0;
00401454
00401467          if (!sub_401364(&buf))
00401484              puts("Wrong!");
00401467          else
00401473              puts("Correct!");
00401473
00401489              result = 0;
00401435      }
00401435      else
00401437          result = 1;
00401437
00401492      *(uint64_t*)((char*)fsbase + 0x28);
00401492
0040149b      if (rax == *(uint64_t*)((char*)fsbase + 0x28))
004014a3          return result;
004014a3
0040149d      __stack_chk_fail();
0040149d      /* no return */
004013eb  }
```

But functions that are part of this program won't have a name associated with them

0x0004010a0 .pl  
0x0004010b0 .pl  
0x0004010c0 .pl  
0x0004010d0 .pl  
0x0004010e0 .pl  
0x0004010f0 .pl  
0x000401100 .pl  
0x000401110 .pl  
0x000401120 .pl  
0x000401130 .pl  
0x000401140 .pl  
0x000401150 .pl  
0x000401160 .te  
0x000401190 .te  
0x0004011c0 .te  
0x000401200 .te  
0x000401240 .te  
0x000401249 .te  
0x0004012db .te  
0x000401364 .te  
0x0004013eb .te  
0x0004014a4 .fi  
0x000403f90 .go  
0x000403f98 .go  
0x000403fa0 .no

es {1}  
es {1}  
if (!sub\_401364(&buf))

```
004013eb int32_t main(int32_t argc, char** argv, char** envp)
{
    void* fsbase;
    int64_t rax = *(uint64_t*)((char*)fsbase + 0x28);
    printf("Enter password: ");
    char buf[0x48];
    int32_t result;

    if (fgets(&buf, 0x40, stdin))
    {
        buf[strcspn(&buf, "\n")] = 0;

        if (!sub_401364(&buf))
            puts("Wrong!");
        else
            puts("Correct!");

        result = 0;
    }
    else
        result = 1;

    *(uint64_t*)((char*)fsbase + 0x28);

    if (rax == *(uint64_t*)((char*)fsbase + 0x28))
        return result;

    __stack_chk_fail();
    /* no return */
}
```

We need to figure out what this function does, and give it a name ourselves :)

```
00401364      uint64_t sub_401364(char* arg1)
```

```
00401364      {
00401364          void* fsbase;
00401374          int64_t rax = *(uint64_t*)((char*)fsbase + 0x28);
00401396          char var_58[0x3f];
00401396          strncpy(&var_58, arg1, 0x40);
0040139b          char var_19 = 0;
004013a6          sub_401249(&var_58);
004013b2          sub_4012db(&var_58);
004013cf          int32_t rax_1;
004013cf          (uint8_t)rax_1 = !strcmp(&var_58, "tfds{zpv_hpu_ju}");
004013d9          *(uint64_t*)((char*)fsbase + 0x28);
004013d9
004013e2          if (rax == *(uint64_t*)((char*)fsbase + 0x28))
004013ea              return (uint64_t)(uint8_t)rax_1;
004013ea
004013e4          __stack_chk_fail();
004013e4          /* no return */
00401364      }
```

```
00401364      uint64_t sub_401364(char* arg1)
```

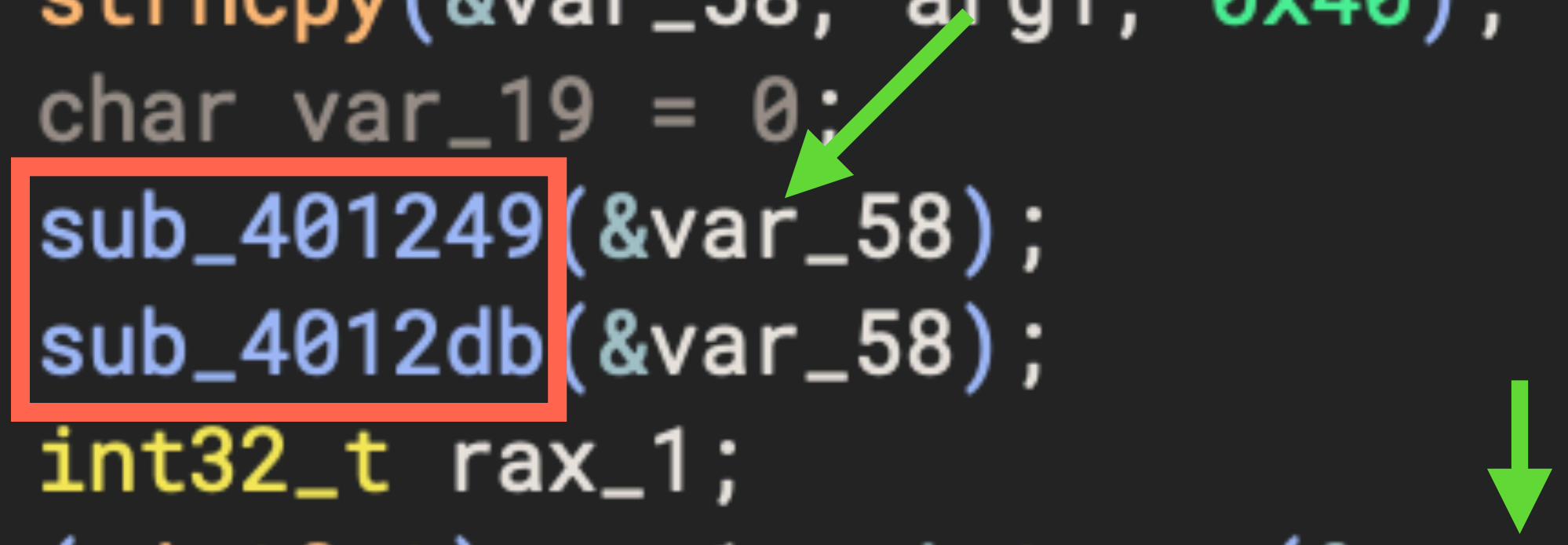
```
00401364      {
00401364          void* fsbase;
00401374          int64_t rax = *(uint64_t*)((char*)fsbase + 0x28);
00401396          char var_58[0x3f];
00401396          strncpy(&var_58, arg1, 0x40);
0040139b          char var_19 = 0;
004013a6          sub_401249(&var_58);
004013b2          sub_4012db(&var_58);
004013cf          int32_t rax_1;
004013cf          (uint8_t)rax_1 = !strcmp(&var_58, "tfds{zpv_hpu_ju}");
004013d9          *(uint64_t*)((char*)fsbase + 0x28);
004013d9
004013e2          if (rax == *(uint64_t*)((char*)fsbase + 0x28))
004013ea              return (uint64_t)(uint8_t)rax_1;
004013ea
004013e4          __stack_chk_fail();
004013e4          /* no return */
00401364      }
```

```
00401364      uint64_t sub_401364(char* arg1)
```

```
00401364      {
00401364          void* fsbase;
00401374          int64_t rax = *(uint64_t*)((char*)fsbase + 0x28);
00401396          char var_58[0x3f];
00401396          strncpy(&var_58, arg1, 0x40);
0040139b          char var_19 = 0;
004013a6          sub_401249(&var_58);
004013b2          sub_4012db(&var_58);
004013cf          int32_t rax_1;
004013cf          (uint8_t)rax_1 = !strcmp(&var_58, "tfds{zpv_hpu_ju}");
004013d9          *(uint64_t*)((char*)fsbase + 0x28);
004013d9
004013e2          if (rax == *(uint64_t*)((char*)fsbase + 0x28))
004013ea              return (uint64_t)(uint8_t)rax_1;
004013ea
004013e4          __stack_chk_fail();
004013e4          /* no return */
00401364      }
```

```
00401364      uint64_t sub_401364(char* arg1)
```

```
00401364      {
00401364          void* fsbase;
00401374          int64_t rax = *(uint64_t*)((char*)fsbase + 0x28);
00401396          char var_58[0x3f];
00401396          strncpy(&var_58, arg1, 0x40);
0040139b          char var_19 = 0;
004013a6          sub_401249(&var_58);
004013b2          sub_4012db(&var_58);
004013cf          int32_t rax_1;
004013cf          (uint8_t)rax_1 = !strcmp(&var_58, "tfds{zpv_hpu_ju}");
004013d9          *(uint64_t*)((char*)fsbase + 0x28);
004013d9
004013e2          if (rax == *(uint64_t*)((char*)fsbase + 0x28))
004013ea              return (uint64_t)(uint8_t)rax_1;
004013ea
004013e4          __stack_chk_fail();
004013e4          /* no return */
00401364      }
```



00401364      uint64\_t sub\_401364(char\* arg1)

```
00401364      {
00401364          void* fsbase;
00401374          int64_t rax = *(uint64_t*)((char*)fsbase + 0x28);
00401396          char var_58[0x3f];
00401396          strncpy(&var_58, arg1, 0x40);      Copy arg1 to some buffer
0040139b          char var_19 = 0;
004013a6          sub_401249(&var_58);      Do some stuff on that buffer
004013b2          sub_4012db(&var_58);
004013cf          int32_t rax_1;
004013cf          (uint8_t)rax_1 = !strcmp(&var_58, "tfds{zpv_hpu_ju}");
004013d9          *(uint64_t*)((char*)fsbase + 0x28);      See if it matches
004013d9                                                      this constant
004013e2          if (rax == *(uint64_t*)((char*)fsbase + 0x28))
004013ea          |      return (uint64_t)(uint8_t)rax_1;
004013ea
004013e4          __stack_chk_fail();
004013e4          /* no return */
00401364      }
```

00401364      uint64\_t sub\_401364(char\* arg1)

```
00401364      {
00401364          void* fsbase;
00401374          int64_t rax = *(uint64_t*)((char*)fsbase + 0x28);
00401396          char var_58[0x3f];
00401396          strncpy(&var_58, arg1, 0x40);      Copy arg1 to some buffer
0040139b          char var_19 = 0;
004013a6          sub_401249(&var_58);      Do some stuff on that buffer
004013b2          sub_4012db(&var_58);
004013cf          int32_t rax_1;
004013cf          (uint8_t)rax_1 = !strcmp(&var_58, "tfds{zpv_hpu_ju}");
004013d9          *(uint64_t*)((char*)fsbase + 0x28);      See if it matches
004013d9                                                      this constant
004013e2          if (rax == *(uint64_t*)((char*)fsbase + 0x28))
004013ea          |      return (uint64_t)(uint8_t)rax_1;
004013ea
004013e4          __stack_chk_fail();
004013e4          /* no return */
00401364      }
```

← → pw\_s × +

# Symbols 🔍 Search symbols

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Name	Address	Sec
sub_401080	0x000401080	.pl
sub_401090	0x000401090	.pl
sub_4010a0	0x0004010a0	.pl
sub_4010b0	0x0004010b0	.pl
__cxa_finalize	0x0004010c0	.pl
strncpy	0x0004010d0	.pl
puts	0x0004010e0	.pl
strlen	0x0004010f0	.pl
__stack_chk_fail	0x000401100	.pl
printf	0x000401110	.pl
strcspn	0x000401120	.pl
fgets	0x000401130	.pl
strcmp	0x000401140	.pl
__ctype_b_loc	0x000401150	.pl
_start	0x000401160	.te
deregister_tm_clon...	0x000401190	.te
sub_4011c0	0x0004011c0	.te
_FINI_0	0x000401200	.te
_INIT_0	0x000401240	.te
sub_401249	0x000401249	.te
sub_4012db	0x0004012db	.te
sub_401364	0x000401364	.te

Cross References

▶ Filter (4)

▼ Code References {2}

▼ sub\_401249 {2}

|→ 00401271 char rax\_5 = arg1[(int64\_t)var\_10];

|→ 004012d7 nop

▼ Variable References {2}

▼ int32\_t var\_10 {1}

|→ 004012d5 if (var\_10 >= result)

▼ int32\_t result {1}

ELF Linear Pseudo C

int64\_t (\* const)() \_INIT\_0()

00401244     return sub\_4011c0();

00401240     }

00401249     int32\_t sub\_401249(char\* arg1)

00401249     {

00401249         int32\_t rax\_1 = strlen(arg1);

00401268         int32\_t var\_10 = 0;

004012d0         int32\_t result;

004012d0         while (true)

004012d0         {

004012d0             result = (rax\_1 + (rax\_1 >> 0x1f)) >> 1;

004012d5             if (var\_10 >= result)

004012d5             |     break;

004012d5             char rax\_5 = arg1[(int64\_t)var\_10];

0040127e             arg1[(int64\_t)var\_10] =

004012a7                 arg1[(int64\_t)(rax\_1 - 1 - var\_10)];

004012a7             arg1[(int64\_t)(rax\_1 - 1 - var\_10)] = rax\_5;

004012c0             var\_10 += 1;

004012c2         }

004012d0         return result;

004012da     }

00401249     }

004012db     int64\_t sub\_4012db(void\* arg1)

004012db     {

004012db         int32\_t var\_c = 0;

00401359         char result;

00401359     }

linux-x86\_64   0x4012d5-0x4012d7 (0x2 bytes)

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← → pw\_s × +

# Symbols 🔍 Search symbols

ELF ▾ Linear ▾ Pseudo C ▾

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Name	Address	Sec
sub_401080	0x000401080	.pl
sub_401090	0x000401090	.pl
sub_4010a0	0x0004010a0	.pl
sub_4010b0	0x0004010b0	.pl
__cxa_finalize	0x0004010c0	.pl
strncpy	0x0004010d0	.pl
puts	0x0004010e0	.pl
strlen	0x0004010f0	.pl
__stack_chk_fail	0x000401100	.pl
printf	0x000401110	.pl
strcspn	0x000401120	.pl
fgets	0x000401130	.pl
strcmp	0x000401140	.pl
__ctype_b_loc	0x000401150	.pl
_start	0x000401160	.pl
deregister_tm_clon...	0x000401190	.pl
sub_4011c0	0x0004011c0	.pl
_FINI_0	0x000401200	.pl
_INIT_0	0x000401240	.pl
sub_401249	0x000401249	.te
sub_4012db	0x0004012db	.te
sub_401364	0x000401364	.te

Cross References ⭐

▶ Filter (4)

▼ Code References {2}

▼ sub\_401249 {2}

|→ 00401271 char rax\_5 = arg1[(int64\_t)var\_10];

|→ 004012d7 nop

▼ Variable References {2}

▼ int32\_t var\_10 {1}

|→ 004012d5 if (var\_10 s>= result)

▼ int32\_t result {1}

int64\_t (\* const)() \_INIT\_0()

00401244     return sub\_4011c0();

00401240     }

00401249     int32\_t sub\_401249(char\* arg1)

00401249     {

00401249         int32\_t rax\_1 = strlen(arg1);

00401268         int32\_t var\_10 = 0;

004012d0         int32\_t result;

004012a7         arg1[(int64\_t)var\_10] =

004012a7             arg1[(int64\_t)(rax\_1 - 1 - var\_10)];

004012c0         arg1[(int64\_t)(rax\_1 - 1 - var\_10)] = rax\_5;

004012c2         var\_10 += 1;

004012d0         }

004012d0         return result;

004012da     }

00401249     }

004012db     int64\_t sub\_4012db(void\* arg1)

004012db     {

004012db         int32\_t var\_c = 0;

00401359         char result;

00401359     }

Define Name

Enter variable name:

rax\_1

Close Accept

linux-x86\_64   0x4012d5-0x4012d7 (0x2 bytes) 🔒

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pw\_s

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Symbols

Q Search symbols

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Name	Address	Sec
sub_401080	0x000401080	.pl
sub_401090	0x000401090	.pl
sub_4010a0	0x0004010a0	.pl
sub_4010b0	0x0004010b0	.pl
__cxa_finalize	0x0004010c0	.pl
strncpy	0x0004010d0	.pl
puts	0x0004010e0	.pl
strlen	0x0004010f0	.pl
__stack_chk_fail	0x000401100	.pl
printf	0x000401110	.pl
strcspn	0x000401120	.pl
fgets	0x000401130	.pl
strcmp	0x000401140	.pl
__ctype_b_loc	0x000401150	.pl
_start	0x000401160	.te
deregister_tm_clon...	0x000401190	.te
sub_4011c0	0x0004011c0	.te
_FINI_0	0x000401200	.te
_INIT_0	0x000401240	.te
sub_401249	0x000401249	.te
sub_4012db	0x0004012db	.te
sub_401364	0x000401364	.te

Cross References

Filter (1)

Code References {1}

sub\_401364 {1}

← 004013b2 sub\_4012db(&password)

ELF Linear Pseudo C

int32\_t sub\_401249(char\* password)

00401244 return sub\_4011c0();

00401240 }

00401249 int32\_t sub\_401249(char\* password)

00401249 {

00401249 int32\_t len = strlen(password);

00401268 int32\_t i = 0;

004012d0 int32\_t result;

004012d0 while (true)

004012d0 {

004012d0 result = (len + (len >> 0x1f)) >> 1;

004012d5 if (i >= result)

004012d5 break;

004012d5 char character = password[i];

0040127e password[i] = password[len - 1 - i];

004012a7 password[len - 1 - i] = character;

004012c0 i += 1;

004012c2 }

004012d0 return result;

004012da }

00401249 }

004012db int64\_t sub\_4012db(void\* arg1)

004012db {

00401359 int32\_t var\_c = 0;

00401359 char result;

00401359 while (true)

linux-x86\_64

0x4012db-0x4012df (0x4 bytes)





← → pw\_s +

# Symbols

Q Search symbols

Name

Address

Sec

sub\_4010800x000401080.pl

sub\_4010900x000401090.pl

sub\_4010a00x0004010a0.pl

sub\_4010b00x0004010b0.pl

\_\_cxa\_finalize0x0004010c0.pl

strncpy0x0004010d0.pl

puts0x0004010e0.pl

strlen0x0004010f0.pl

\_\_stack\_chk\_fail0x000401100.pl

printf0x000401110.pl

strcspn0x000401120.pl

fgets0x000401130.pl

strcmp0x000401140.pl

\_\_ctype\_b\_loc0x000401150.pl

\_start0x000401160.te

deregister\_tm\_clon...0x000401190.te

sub\_4011c00x0004011c0.te

\_FINI\_00x000401200.te

\_INIT\_00x000401240.te

sub\_4012490x000401249.te

sub\_4012db0x0004012db.te

sub\_4013640x000401364.te

Cross References

Filter (1)

Code References {1}

sub\_401364 {1}

← 004013b2 sub\_4012db(&password)

ELF Linear Pseudo C

int32\_t sub\_401249(char\* password)

00401244return sub\_4011c0();

00401240}

00401249int32\_t sub\_401249(char\* password)

00401249{

00401249int32\_t len = strlen(password);

00401268int32\_t i = 0;

004012d0int32\_t result;

004012d0while (true)

004012d0{

004012d0result = (len + (len >> 0x1f)) >> 1;

004012d0

004012d5if (i >= result)

004012d5break;

004012d5

0040127echar character = password[i];

004012a7password[i] = password[len - 1 - i];

004012c0password[len - 1 - i] = character;

004012c2i += 1;

004012d0}

004012d0return result;

004012da

00401249}

004012dbint64\_t sub\_4012db(void\* arg1)

004012db{

00401359int32\_t var\_c = 0;

00401359char result;

00401359while (true)

linux-x86\_64 0x4012db-0x4012df (0x4 bytes)

Optimised signed integer halving



← → pw\_s +

# Symbols 🔍 Search symbols

Name	Address	Sec
sub_401080	0x000401080	.pl
sub_401090	0x000401090	.pl
sub_4010a0	0x0004010a0	.pl
sub_4010b0	0x0004010b0	.pl
__cxa_finalize	0x0004010c0	.pl
strncpy	0x0004010d0	.pl
puts	0x0004010e0	.pl
strlen	0x0004010f0	.pl
__stack_chk_fail	0x000401100	.pl
printf	0x000401110	.pl
strcspn	0x000401120	.pl
fgets	0x000401130	.pl
strcmp	0x000401140	.pl
__ctype_b_loc	0x000401150	.pl
_start	0x000401160	.te
deregister_tm_clon...	0x000401190	.te
sub_4011c0	0x0004011c0	.te
_FINI_0	0x000401200	.te
_INIT_0	0x000401240	.te
sub_401249	0x000401249	.te
sub_4012db	0x0004012db	.te
sub_401364	0x000401364	.te

Cross References

Filter (1)

Code References {1}

sub\_401364 {1}

← 004013b2 sub\_4012db(&password)

ELF Linear Pseudo C

int32\_t sub\_401249(char\* password)

00401244 return sub\_4011c0();

00401240 }

00401249 int32\_t sub\_401249(char\* password)

00401249 {

00401249 int32\_t len = strlen(password);

00401268 int32\_t i = 0;

004012d0 int32\_t result;

004012d0 while (true)

004012d0 {

004012d0 result = (len + (len >> 0x1f)) >> 1;

004012d0

004012d5 if (i >= result)

004012d5 break;

004012d5

004012e char character = password[i];

004012a7 password[i] = password[len - 1 - i];

004012c0 password[len - 1 - i] = character;

004012c2 i += 1;

004012d0 }

004012d0 return result;

004012da

00401249 }

004012db int64\_t sub\_4012db(void\* arg1)

004012db {

00401359 int32\_t var\_c = 0;

00401359 char result;

00401359 while (true)

Optimised signed integer halving

Go test things out in python if you're unsure

linux-x86\_64 0x4012db-0x4012df (0x4 bytes)

← → pw\_s +

# Symbols 🔍 Search symbols

ELF ▾ Linear ▾ Pseudo C ▾

🔗 📄 ☰

Name	Address	Sec
sub_401080	0x000401080	.pl
sub_401090	0x000401090	.pl
sub_4010a0	0x0004010a0	.pl
sub_4010b0	0x0004010b0	.pl
__cxa_finalize	0x0004010c0	.pl
strncpy	0x0004010d0	.pl
puts	0x0004010e0	.pl
strlen	0x0004010f0	.pl
__stack_chk_fail	0x000401100	.pl
printf	0x000401110	.pl
strcspn	0x000401120	.pl
fgets	0x000401130	.pl
strcmp	0x000401140	.pl
__ctype_b_loc	0x000401150	.pl
_start	0x000401160	.te
deregister_tm_clon...	0x000401190	.te
sub_4011c0	0x0004011c0	.te
_FINI_0	0x000401200	.te
_INIT_0	0x000401240	.te
sub_401249	0x000401249	.te
sub_4012db	0x0004012db	.te
sub_401364	0x000401364	.te

Cross References ⭐

▶ Filter (1)

▼ Code References {1}

▼ sub\_401364 {1}

← 004013b2 sub\_4012db(&password)

🔄 int32\_t sub\_401249(char\* password)

00401244 | return sub\_4011c0();

00401240 }

00401249 int32\_t sub\_401249(char\* password)

00401249 {

00401249 int32\_t len = strlen(password);

00401268 int32\_t i = 0;

004012d0 int32\_t result;

004012d0 while (true)

004012d0 {

004012d0 result = (len + (len >> 0x1f)) >> 1;

004012d5 if (i >= result)

004012d5 break;

004012d5

0040127e char character = password[i];

004012a7 password[i] = password[len - 1 - i];

004012c0 password[len - 1 - i] = character;

004012c2 i += 1;

004012d0 }

004012d0

004012d0 return result;

004012da

00401249 }

004012db int64\_t sub\_4012db(void\* arg1)

004012db {

00401359 int32\_t var\_c = 0;

00401359 char result;

00401359 while (true)

Loop exit

linux-x86\_64 0x4012db-0x4012df (0x4 bytes)



pw\_s

Symbols Search symbols

Name	Address	Sec
sub_401080	0x000401080	.pl
sub_401090	0x000401090	.pl
sub_4010a0	0x0004010a0	.pl
sub_4010b0	0x0004010b0	.pl
__cxa_finalize	0x0004010c0	.pl
strncpy	0x0004010d0	.pl
puts	0x0004010e0	.pl
strlen	0x0004010f0	.pl
__stack_chk_fail	0x000401100	.pl
printf	0x000401110	.pl
strcspn	0x000401120	.pl
fgets	0x000401130	.pl
strcmp	0x000401140	.pl
__ctype_b_loc	0x000401150	.pl
_start	0x000401160	.te
deregister_tm_clon...	0x000401190	.te
sub_4011c0	0x0004011c0	.te
_FINI_0	0x000401200	.te
_INIT_0	0x000401240	.te
sub_401249	0x000401249	.te
sub_4012db	0x0004012db	.te
sub_401364	0x000401364	.te

Cross References

Filter (1)

Code References {1}

sub\_401364 {1}

← 004013b2 sub\_4012db(&password)

ELF Linear Pseudo C

```
int32_t sub_401249(char* password)
{
    return sub_4011c0();
}

int32_t sub_401249(char* password)
{
    int32_t len = strlen(password);
    int32_t i = 0;
    int32_t result;

    while (true)
    {
        result = (len + (len >> 0x1f)) >> 1;

        if (i >= result)
            break;

        char character = password[i];
        password[i] = password[len - 1 - i];
        password[len - 1 - i] = character;
        i += 1;
    }

    return result;
}

int64_t sub_4012db(void* arg1)
{
    int32_t var_c = 0;
    char result;

    while (true)
```

Swap the  $i$ th character with the  $i$ th last one

linux-x86\_64 0x4012db-0x4012df (0x4 bytes)

← → pw\_s +

# Symbols

Search symbols

Name	Address	Sec
sub_401080	0x000401080	.pl
sub_401090	0x000401090	.pl
sub_4010a0	0x0004010a0	.pl
sub_4010b0	0x0004010b0	.pl
__cxa_finalize	0x0004010c0	.pl
strncpy	0x0004010d0	.pl
puts	0x0004010e0	.pl
strlen	0x0004010f0	.pl
__stack_chk_fail	0x000401100	.pl
printf	0x000401110	.pl
strcspn	0x000401120	.pl
fgets	0x000401130	.pl
strcmp	0x000401140	.pl
__ctype_b_loc	0x000401150	.pl
_start	0x000401160	.te
deregister_tm_clon...	0x000401190	.te
sub_4011c0	0x0004011c0	.te
_FINI_0	0x000401200	.te
_INIT_0	0x000401240	.te
sub_401249	0x000401249	.te
sub_4012db	0x0004012db	.te
sub_401364	0x000401364	.te

Cross References

Filter (1)

Code References {1}

sub\_401364 {1}

← 004013b2 sub\_4012db(&password)

ELF Linear Pseudo C

int32\_t sub\_401249(char\* password)

00401244 return sub\_4011c0();

00401240 }

00401249 int32\_t sub\_401249(char\* password)

00401249 {

00401249 int32\_t len = strlen(password);

00401268 int32\_t i = 0;

004012d0 int32\_t result;

004012d0 while (true)

004012d0 {

004012d0 result = (len + (len >> 0x1f)) >> 1;

004012d5 if (i >= result)

004012d5 break;

004012d5

0040127e char character = password[i];

004012a7 password[i] = password[len - 1 - i];

004012c0 password[len - 1 - i] = character;

004012c2 i += 1;

004012d0 }

004012da return result;

00401249 }

004012db int64\_t sub\_4012db(void\* arg1)

004012db {

00401359 int32\_t var\_c = 0;

00401359 char result;

00401359 while (true)

linux-x86\_64 0x4012db-0x4012df (0x4 bytes)

string reversal!



00401364 uint64\_t sub\_401364(char\* arg1)

```
00401364 {
00401364     void* fsbase;
00401374     int64_t rax = *(fsbase + 0x28);
00401396     char password[0x3f];
00401396     strncpy(&password, arg1, 0x40);
0040139b     char var_19 = 0;
004013a6     reverse_string(&password);
004013b2     sub_4012db(&password);
004013cf     int32_t rax_1;
004013cf     rax_1 = !strcmp(&password, "tfds{zpv_hpu_ju}");
004013d9     *(fsbase + 0x28);
004013d9
004013e2     if (rax == *(fsbase + 0x28))
004013ea     |         return rax_1;
004013ea
004013e4     __stack_chk_fail();
004013e4     /* no return */
00401364 }
```

004012db int64\_t sub\_4012db(void\* arg1)

```
004012db {
004012db     int32_t var_c = 0;
00401359     char result;
00401359
00401359     while (true)
00401359     {
00401359         result = *(arg1 + var_c);
00401359
00401359         if (!result)
0040135e             break;
0040135e
0040135e         if ((*__ctype_b_loc())[*(arg1 + var_c)] & 0x400)
00401322             *(arg1 + var_c) += 1;
00401346
00401346         var_c += 1;
00401348
00401359     }
00401359
00401359     return result;
00401363
004012db }
```





```
004012db void sub_4012db(char* password)
```

```
004012db {
004012db     int32_t i = 0;
004012db
0040135e while (password[i])
0040135e {
0040135e     if ((*__ctype_b_loc())[password[i]] & 0x400)
00401346         password[i] += 1;
00401346
00401348     i += 1;
0040135e }
004012db }
```

004012db

```
void sub_4012db(char* password)
```

004012db

```
{
```

004012db

```
    int32_t i = 0;
```

004012db

0040135e

```
    while (password[i])
```

0040135e

```
{
```

0040135e

```
        if ((*__ctype_b_loc())[password[i]] & 0x400)
```

00401346

```
            password[i] += 1;
```

00401346

00401348

```
            i += 1;
```

0040135e

```
        }
```

004012db

```
}
```

**Checks if the character is alphanumeric**  
in C this would be `isAlpha`

004012db void sub\_4012db(char\* password)

```
004012db {
004012db     int32_t i = 0;
004012db
0040135e while (password[i])
0040135e {
0040135e     if ((*__ctype_b_loc())[password[i]] & 0x400)
00401346         password[i] += 1;
00401346
00401348     i += 1;
0040135e }
004012db }
```

Shifts every character by 1!

```
bool validate_password(char* pw_inp)
```

```
{  
    void* fsbase;  
    int64_t rax = *(fsbase + 0x28);  
    char pw_buf[0x3f];  
    strncpy(&pw_buf, pw_inp, 0x40);  
    char var_19 = 0;  
    reverse_string(&pw_buf);  
    shift_chars(&pw_buf);  
    bool is_valid = !strcmp(&pw_buf, "tfds{zpv_hpu_ju}");  
    *(fsbase + 0x28);  
  
    if (rax == *(fsbase + 0x28))  
        return is_valid;  
  
    __stack_chk_fail();  
    /* no return */  
}
```



00401364 uint64\_t sub\_401364(char\* arg1)

```
00401364 {
00401364     void* fsbase;
00401374     int64_t rax = *(uint64_t*)((char*)fsbase + 0x28);
00401396     char var_58[0x3f];
00401396     strncpy(&var_58, arg1, 0x40);
0040139b     char var_19 = 0;
004013a6     sub_401249(&var_58);
004013b2     sub_4012db(&var_58);
004013cf     int32_t rax_1;
004013cf     (uint8_t)rax_1 = !strcmp(&var_58, "tfds{zpv_hpu_ju}");
004013d9     *(uint64_t*)((char*)fsbase + 0x28);
004013d9
004013e2     if (rax == *(uint64_t*)((char*)fsbase + 0x28))
004013ea     |         return (uint64_t)(uint8_t)rax_1;
004013ea
004013e4     __stack_chk_fail();
004013e4     /* no return */
00401364 }
```



```
bool validate_password(char* pw_inp)
```

```
{  
    void* fsbase;  
    int64_t rax = *(fsbase + 0x28);  
    char pw_buf[0x3f];  
    strncpy(&pw_buf, pw_inp, 0x40);  
    char var_19 = 0;  
    reverse_string(&pw_buf);  
    shift_chars(&pw_buf);  
    bool is_valid = !strcmp(&pw_buf, "tfds{zpv_hpu_ju}");  
    *(fsbase + 0x28);  
  
    if (rax == *(fsbase + 0x28))  
        return is_valid;  
  
    __stack_chk_fail();  
    /* no return */  
}
```





```
bool validate_password(char* pw_inp)
{
    void* fsbase;
    int64_t rax = *(fsbase + 0x28);
    char pw_buf[0x3f];
    strncpy(&pw_buf, pw_inp, 0x40);
    char var_19 = 0;
    reverse_string(&pw_buf);
    shift_chars(&pw_buf);
    bool is_valid = !strcmp(&pw_buf, "tfds{zpv_hpu_ju}");
    *(fsbase + 0x28);

    if (rax == *(fsbase + 0x28))
        return is_valid;

    __stack_chk_fail();
    /* no return */
}
```

```
int check_password(const char *input) {
    char buf[64];
    strncpy(buf, input, sizeof(buf));
    buf[sizeof(buf)-1] = 0;

    reverse(buf);
    shift(buf);

    return strcmp(buf, "tfds{zpv_hpu_ju}") == 0;
}
```



```
bool validate_password(char* pw_inp)
{
    void* fsbase;
    int64_t rax = *(fsbase + 0x28);
    char pw_buf[0x3f];
    strncpy(&pw_buf, pw_inp, 0x40);
    char var_19 = 0;
    reverse_string(&pw_buf);
    shift_chars(&pw_buf);
    bool is_valid = !strcmp(&pw_buf, "tfds{zpv_hpu_ju}");
    *(fsbase + 0x28);

    if (rax == *(fsbase + 0x28))
        return is_valid;

    __stack_chk_fail();
    /* no return */
}
```

```
int check_password(const char *input) {
    char buf[64];
    strncpy(buf, input, sizeof(buf));
    buf[sizeof(buf)-1] = 0;

    reverse(buf);
    shift(buf);

    return strcmp(buf, "tfds{zpv_hpu_ju}") == 0;
}
```

***Decompilation is not 1:1***

# Memory...

... can be used as a stack



- From before: `mov DWORD PTR [rbp-0x4], edi`

# Memory...

... can be used as a stack



- From before: `mov DWORD PTR [rbp-0x4], edi`
- What's going on here?

# Memory...

... can be used as a stack



- From before: `mov DWORD PTR [rbp-0x4], edi`
- What's going on here?
- This is putting an item (`edi`) into some location in the stack (`rbp-0x4`)



# Memory...

... can be used as a stack

- From before: `mov DWORD PTR [rbp-0x4], edi`
- What's going on here?
- This is putting an item (`edi`) into some location in the stack (`rbp-0x4`)
- What is a stack?



# Memory...

... can be used as a stack

- From before: `mov DWORD PTR [rbp-0x4], edi`
- What's going on here?
- This is putting an item (`edi`) into some location in the stack (`rbp-0x4`)
- What is a stack?
- What is `rbp`?



# Memory...

... can be used as a stack

- From before: `mov DWORD PTR [rbp-0x4], edi`
- What's going on here?
- This is putting an item (`edi`) into some location in the stack (`rbp-0x4`)
- What is a stack?
- What is `rbp`?
- Why haven't I mentioned `rsp`?

# Memory...

... can be used as a stack

- The "stack" is a section of memory used during execution



# Memory...

... can be used as a stack

- The "stack" is a section of memory used during execution
- Every function allocates it's own stack space on entry



# Memory...

... can be used as a stack



- The "stack" is a section of memory used during execution
- Every function allocates it's own stack space on entry
  - and thus returns it back to it's previous state on exit

# Memory...

... can be used as a stack



- The "stack" is a section of memory used during execution
- Every function allocates it's own stack space on entry
  - and thus returns it back to it's previous state on exit
- Allows for push/pop from the "top" of the stack (much like the ADT)

# Memory...

... can be used as a stack

```
push 0x41
push 0x42
push 0x43
pop  eax
push 0x44
push 0x45
pop  ebx
pop  ecx
push 0x46
push 0x47
push 0x48
pop  edx
pop  esi
pop  edi
pop  ebp
```

stack pointer →

Address	Value
0	
8	
16	
24	
32	
40	



# Memory...

... can be used as a stack

```
push 0x41
push 0x42
push 0x43
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push 0x48
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pop  edi
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```



stack pointer →

Address	Value
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16	
24	
32	
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... can be used as a stack

```
push 0x41
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pop  ebx
pop  ecx
push 0x46
push 0x47
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pop  edi
pop  ebp
```

stack pointer →

Address	Value
0	0x41
8	
16	
24	
32	
40	



# Memory...

... can be used as a stack

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push 0x41
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pop  ebx
pop  ecx
push 0x46
push 0x47
push 0x48
pop  edx
pop  esi
pop  edi
pop  ebp
```

stack pointer →

Address	Value
0	0x41
8	0x42
16	
24	
32	
40	



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... can be used as a stack

```
push 0x41
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pop  esi
pop  edi
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```



stack pointer →

Address	Value
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# Memory...

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pop  esi
pop  edi
pop  ebp
```



stack pointer →

Address	Value
0	0x41
8	0x42
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24	
32	
40	

# Memory...

... can be used as a stack

```
push 0x41
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push 0x47
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pop  edx
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pop  edi
pop  ebp
```

eax: 0x43

stack pointer →

Address	Value
0	0x41
8	0x42
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# Memory...

... can be used as a stack



```
push 0x41
push 0x42
push 0x43
pop  eax
push 0x44
push 0x45
pop  ebx
pop  ecx
push 0x46
push 0x47
push 0x48
pop  edx
pop  esi
pop  edi
pop  ebp
```

eax: 0x43

stack pointer →

Address	Value
0	0x41
8	0x42
16	0x44
24	
32	
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# Memory...

... can be used as a stack



```
push 0x41
push 0x42
push 0x43
pop  eax
push 0x44
push 0x45
pop  ebx
pop  ecx
push 0x46
push 0x47
push 0x48
pop  edx
pop  esi
pop  edi
pop  ebp
```

eax: 0x43

stack pointer →

Address	Value
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32	
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# Memory...

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```
push 0x41
push 0x42
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push 0x44
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pop  ebx
pop  ecx
push 0x46
push 0x47
push 0x48
pop  edx
pop  esi
pop  edi
pop  ebp
```

```
eax: 0x43
ebx: 0x45
```

stack pointer →

Address	Value
0	0x41
8	0x42
16	0x44
24	0x45
32	
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# Memory...

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push 0x41
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pop  ecx
push 0x46
push 0x47
push 0x48
pop  edx
pop  esi
pop  edi
pop  ebp
```

```
eax: 0x43
ebx: 0x45
ecx: 0x44
```

stack pointer →

Address	Value
0	0x41
8	0x42
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```

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ecx: 0x44
```

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```

```
eax: 0x43
ebx: 0x45
ecx: 0x44
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stack pointer →

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```

```
eax: 0x43
ebx: 0x45
ecx: 0x44
```

Address	Value
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stack pointer →

# Memory...

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```

```
eax: 0x43
ebx: 0x45
ecx: 0x44
edx: 0x48
```

stack pointer →

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# Memory...

... can be used as a stack



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pop  ecx
push 0x46
push 0x47
push 0x48
pop  edx
pop  esi
pop  edi
pop  ebp
```

```
eax: 0x43
ebx: 0x45
ecx: 0x44
edx: 0x48
esi: 0x47
```

stack pointer →

Address	Value
0	0x41
8	0x42
16	0x46
24	0x47
32	0x48
40	

# Memory...

... can be used as a stack



```
push 0x41
push 0x42
push 0x43
pop  eax
push 0x44
push 0x45
pop  ebx
pop  ecx
push 0x46
push 0x47
push 0x48
pop  edx
pop  esi
pop  edi
pop  ebp
```

```
eax: 0x43
ebx: 0x45
ecx: 0x44
edx: 0x48
esi: 0x47
edi: 0x48
```

stack pointer →

Address	Value
0	0x41
8	0x42
16	0x46
24	0x47
32	0x48
40	

# Memory...

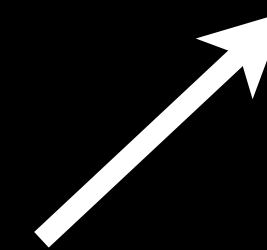
... can be used as a stack



```
push 0x41
push 0x42
push 0x43
pop  eax
push 0x44
push 0x45
pop  ebx
pop  ecx
push 0x46
push 0x47
push 0x48
pop  edx
pop  esi
pop  edi
pop  ebp
```

```
eax: 0x43
ebx: 0x45
ecx: 0x44
edx: 0x48
esi: 0x47
edi: 0x48
ebp: 0x42
```

stack pointer



Address	Value
0	0x41
8	0x42
16	0x46
24	0x47
32	0x48
40	

# Memory...

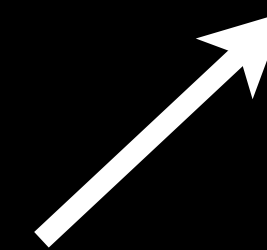
... can be used as a stack



```
push 0x41
push 0x42
push 0x43
pop  eax
push 0x44
push 0x45
pop  ebx
pop  ecx
push 0x46
push 0x47
push 0x48
pop  edx
pop  esi
pop  edi
pop  ebp
```

```
eax: 0x43
ebx: 0x45
ecx: 0x44
edx: 0x48
esi: 0x47
edi: 0x48
ebp: 0x42
```

stack pointer  
esp/rsp



Address	Value
0	0x41
8	0x42
16	0x46
24	0x47
32	0x48
40	

# Memory...

...needs to be function local



- *"Every function allocates it's own stack space on entry"*

# Memory...

...needs to be function local



- *"Every function allocates it's own stack space on entry"*
- How is this accomplished?

# Memory...

...needs to be function local



- *"Every function allocates it's own stack space on entry"*
- How is this accomplished?
- Start of function (prologue) sets up a "frame"

## prologue

```
push    rbp
mov     rbp, rsp
sub     rsp, N
```



# Memory...

...needs to be function local

- *"Every function allocates it's own stack space on entry"*
- How is this accomplished?
- Start of function (prologue) sets up a "frame"
- End of function (epilogue) goes back to the previous (caller's) frame

## prologue

```
push    rbp
mov     rbp, rsp
sub     rsp, N
```

## epilogue

```
mov     rsp, rbp
pop     rbp
ret
```

# Memory...

...needs to be function local

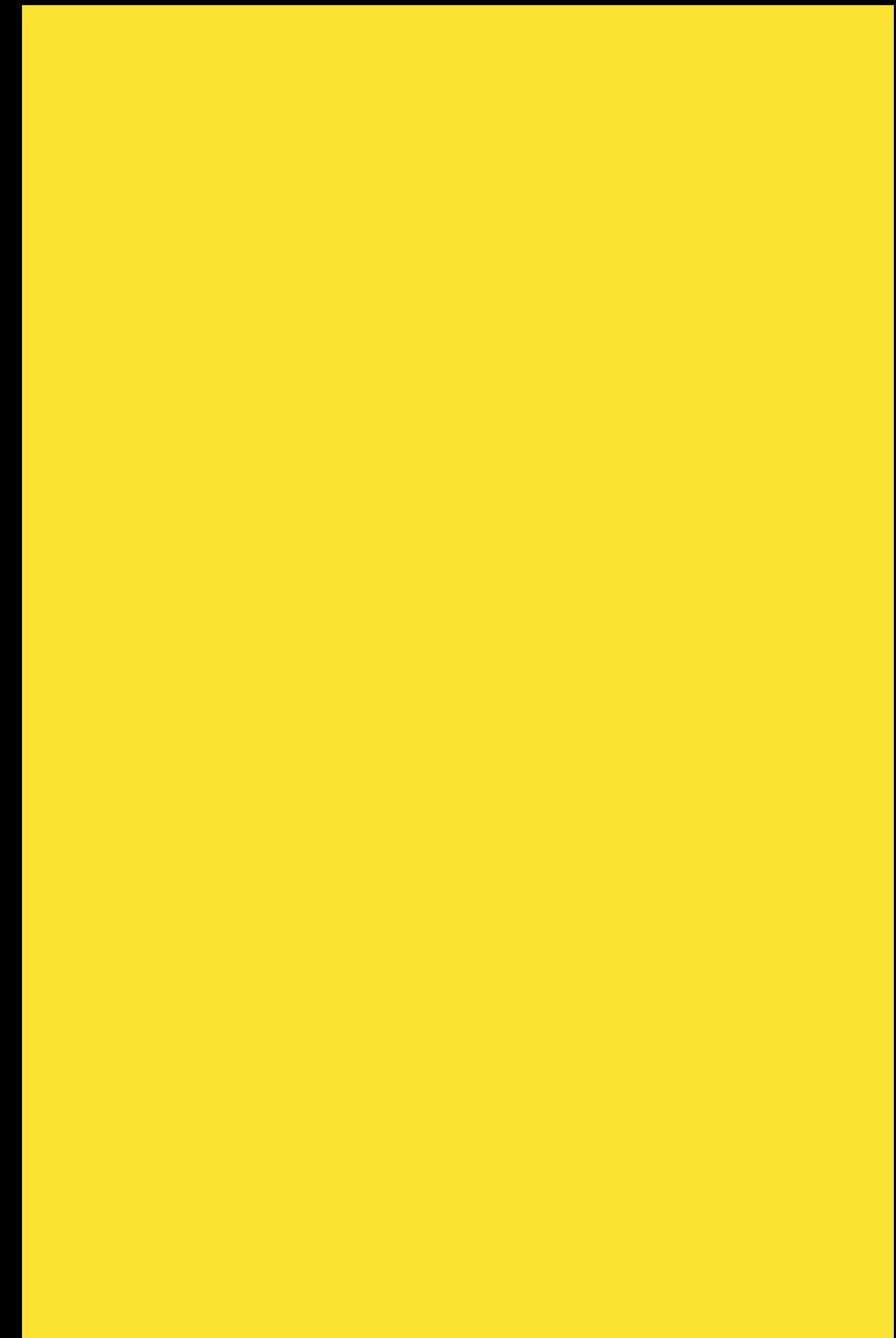


## prologue

```
push    rbp
mov     rbp, rsp
sub     rsp, N
```

rsp →

Stack



# Memory...

...needs to be function local



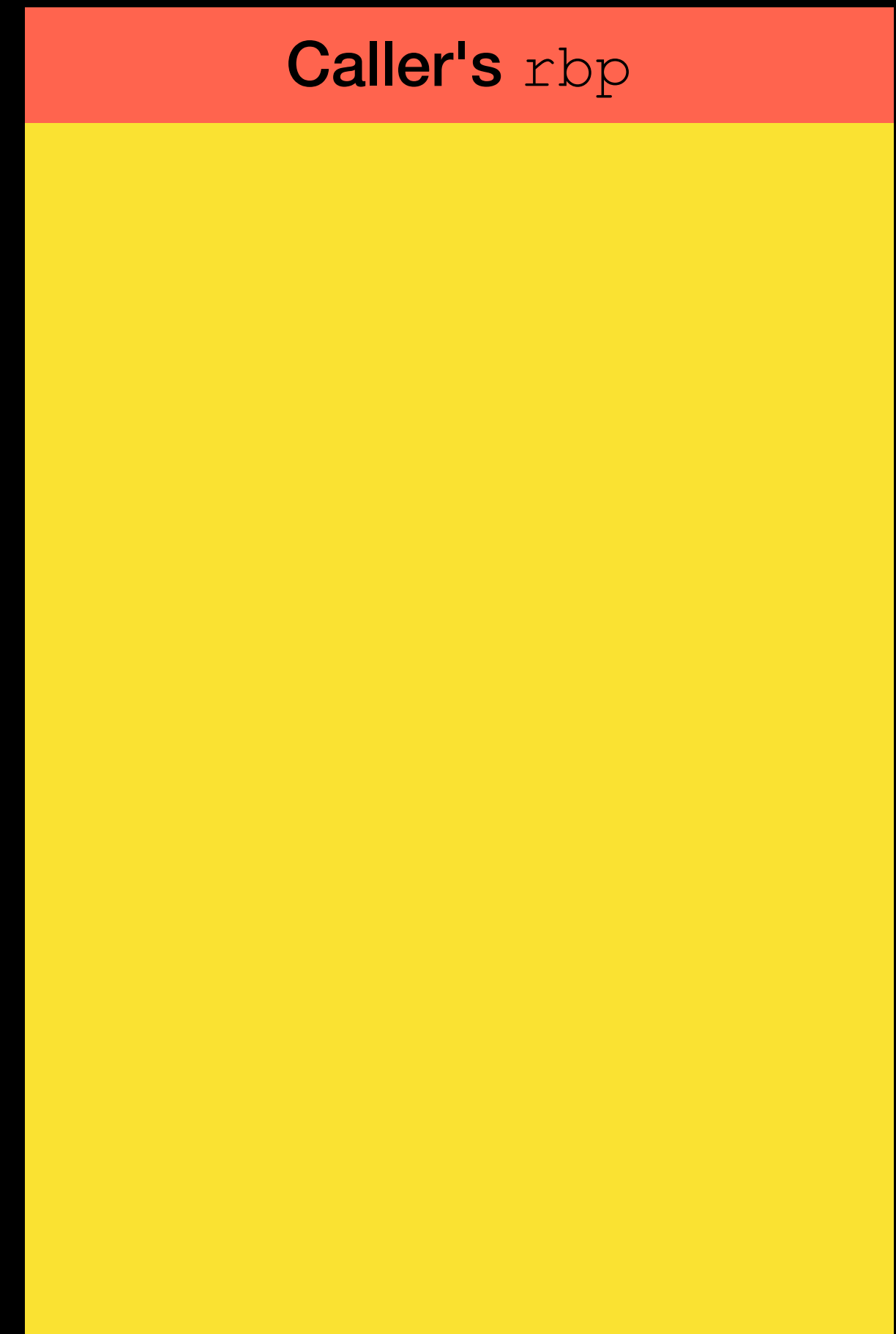
## prologue

```
push    rbp
mov     rbp, rsp
sub     rsp, N
```

rsp →

Stack

Caller's rbp



# Memory...

...needs to be function local



Stack

**prologue**  
push rbp  
mov rbp, rsp  
sub rsp, N

rsp, rbp →



# Memory...

...needs to be function local



## prologue

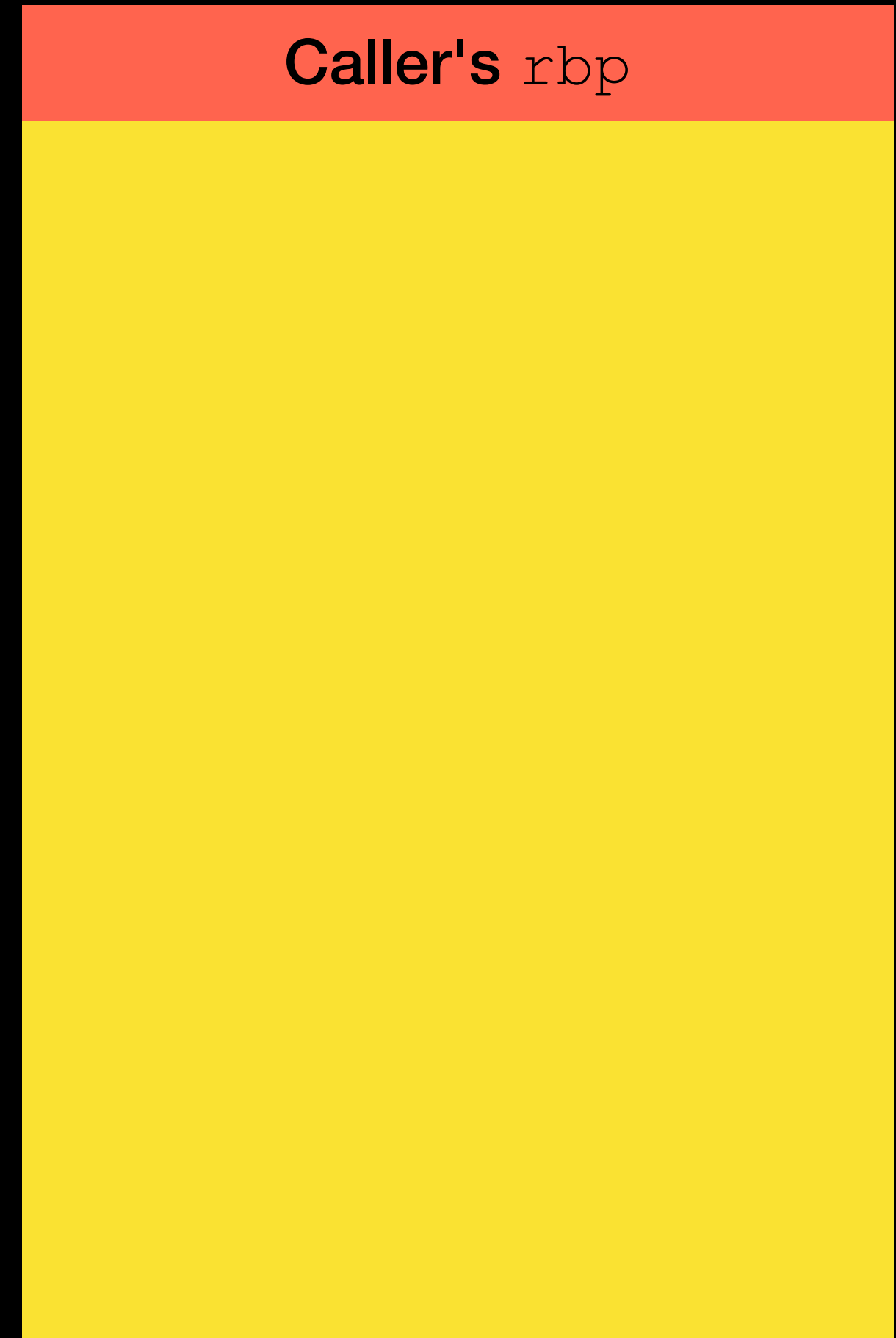
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mov     rbp, rsp
sub     rsp, N
```

Stack

rbp →

Caller's rbp

rsp →



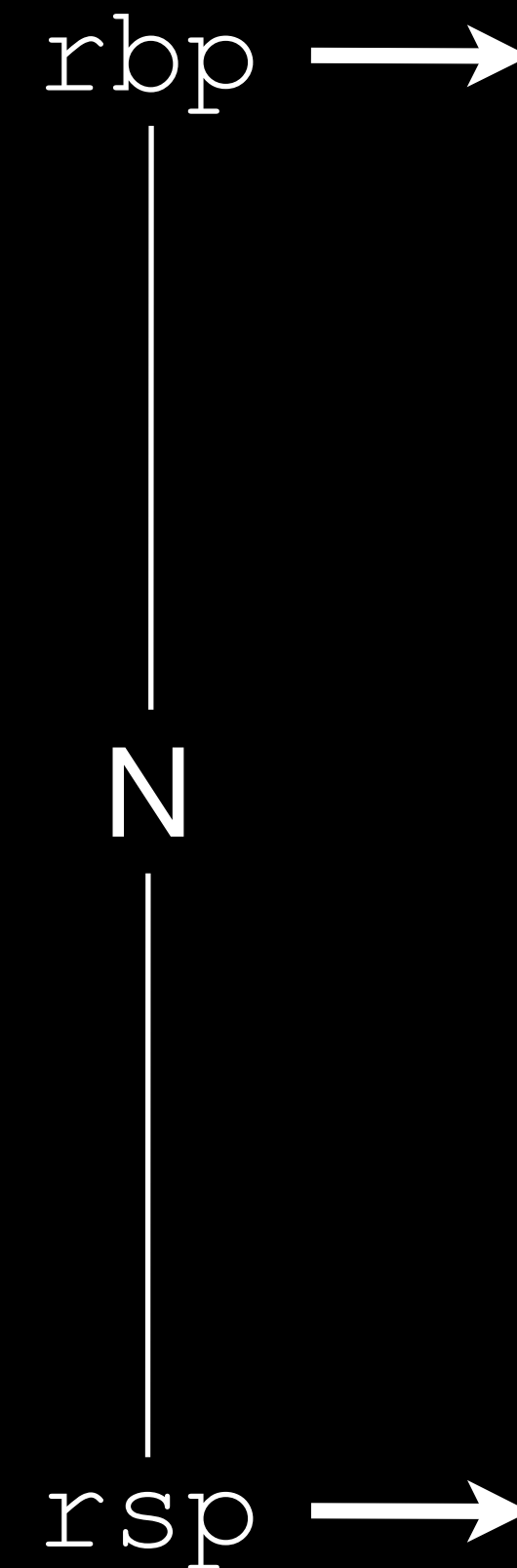
# Memory...

...needs to be function local



Stack

**prologue**  
push rbp  
mov rbp, rsp  
sub rsp, N



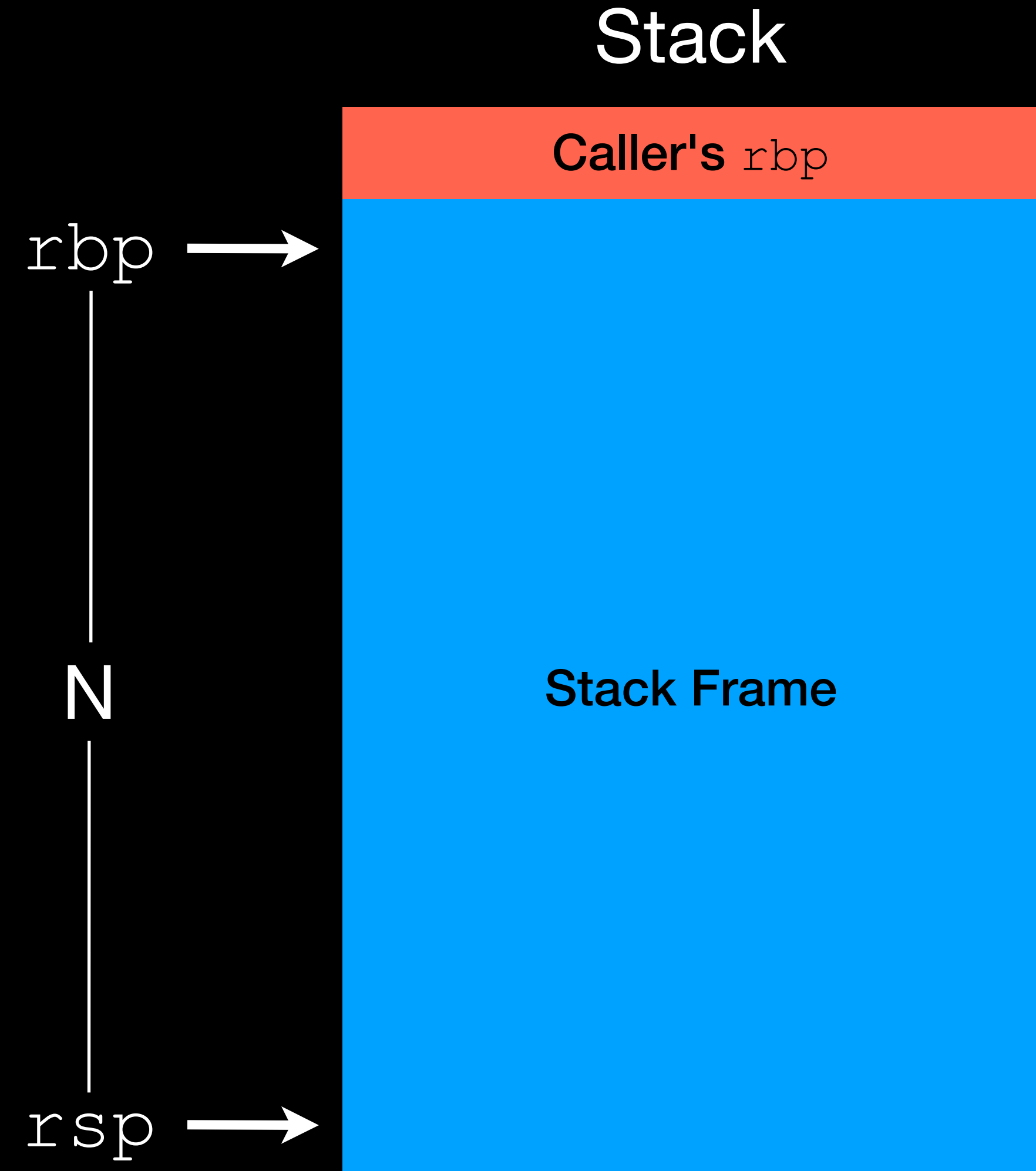
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**prologue**

```
push    rbp
mov     rbp, rsp
sub     rsp, N
```



# Memory...

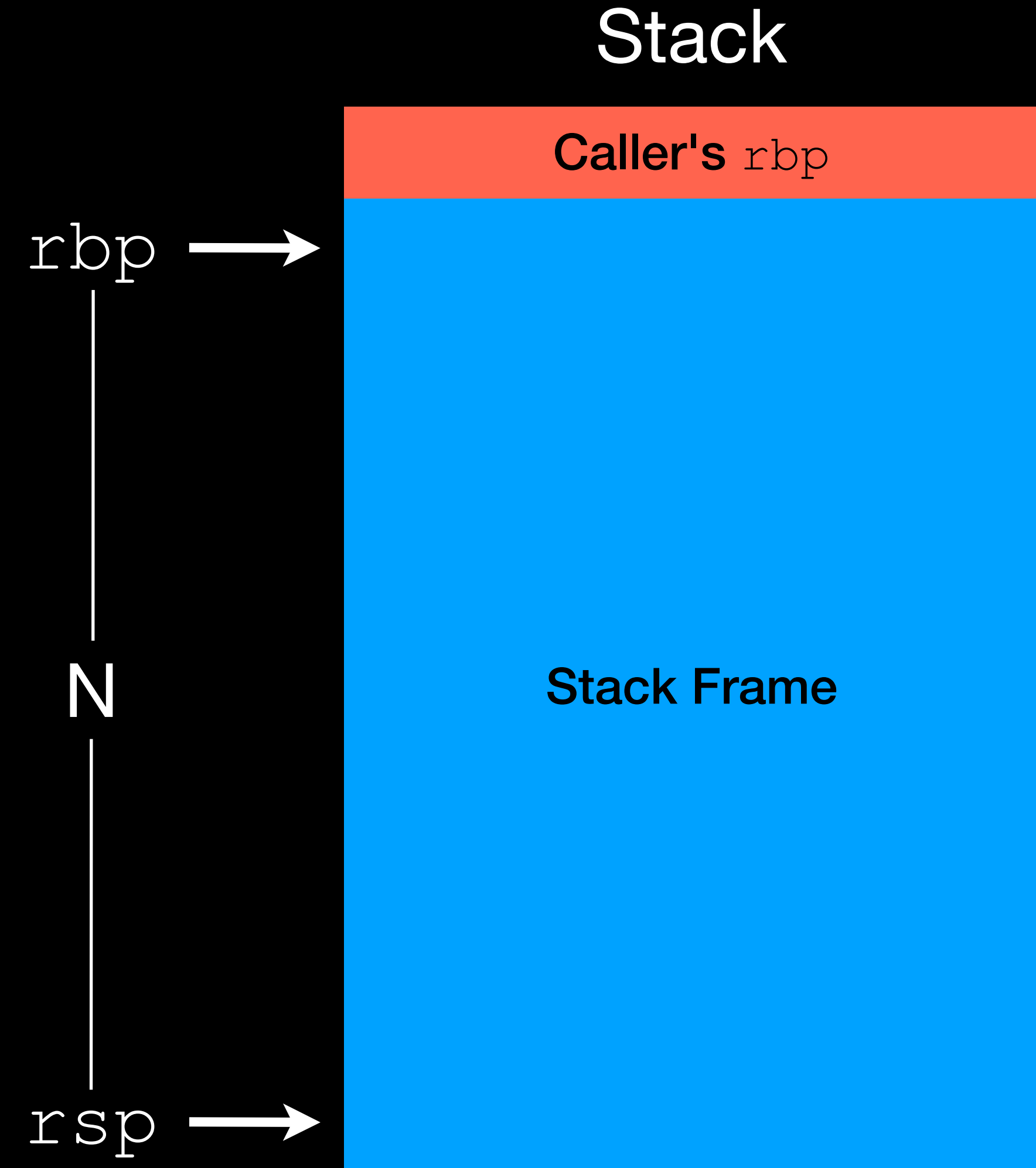
...needs to be function local



## prologue

```
push    rbp
mov     rbp, rsp
sub     rsp, N
```

```
mov     DWORD PTR [rbp-0x4], edi
```



# Memory...

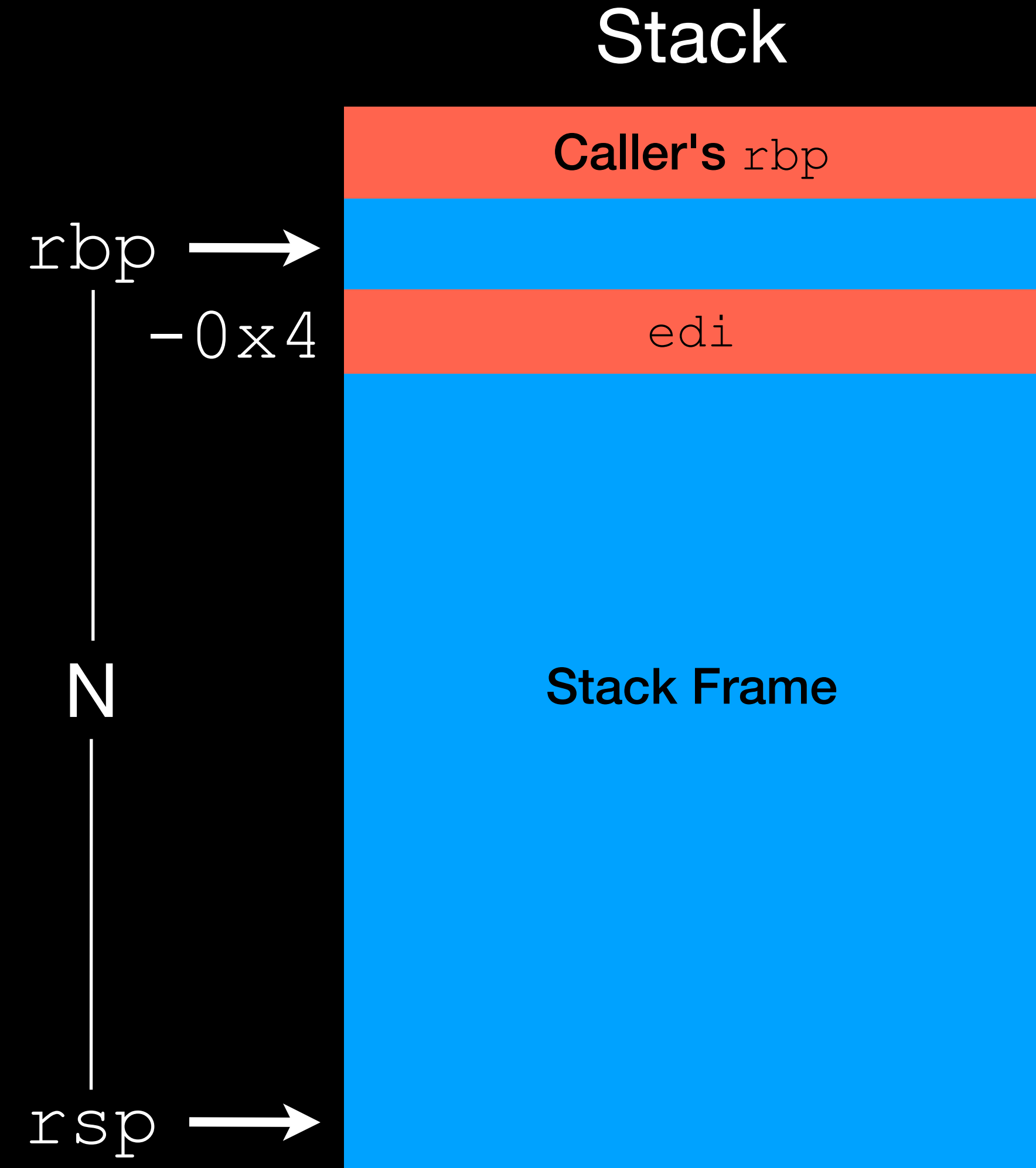
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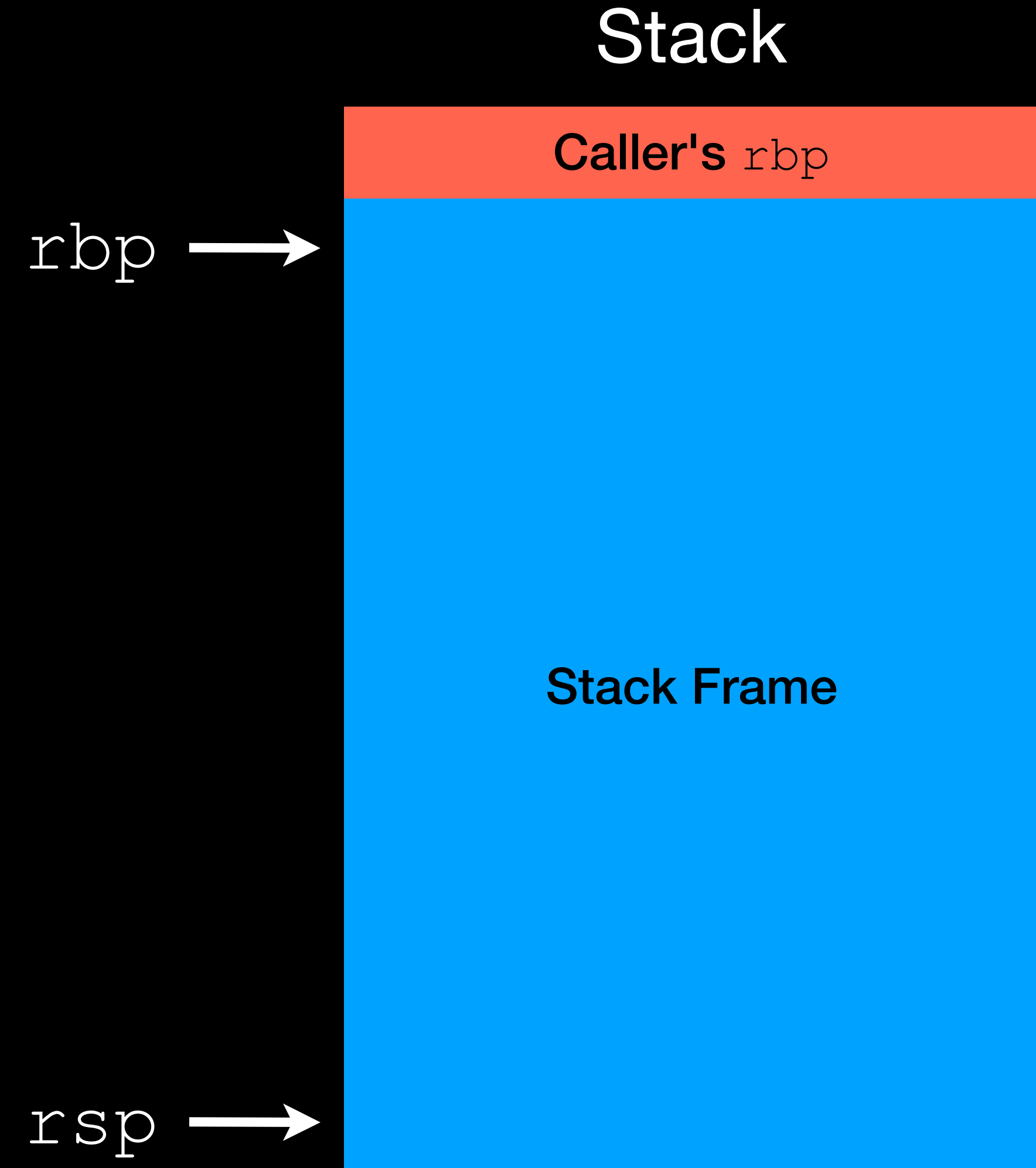
# Memory...

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**epilogue**

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pop    rbp
ret
```



# Memory...

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Stack

**epilogue**

```
mov    rsp, rbp
pop    rbp
ret
```

`rsp, rbp` →



# Memory...

...needs to be function local

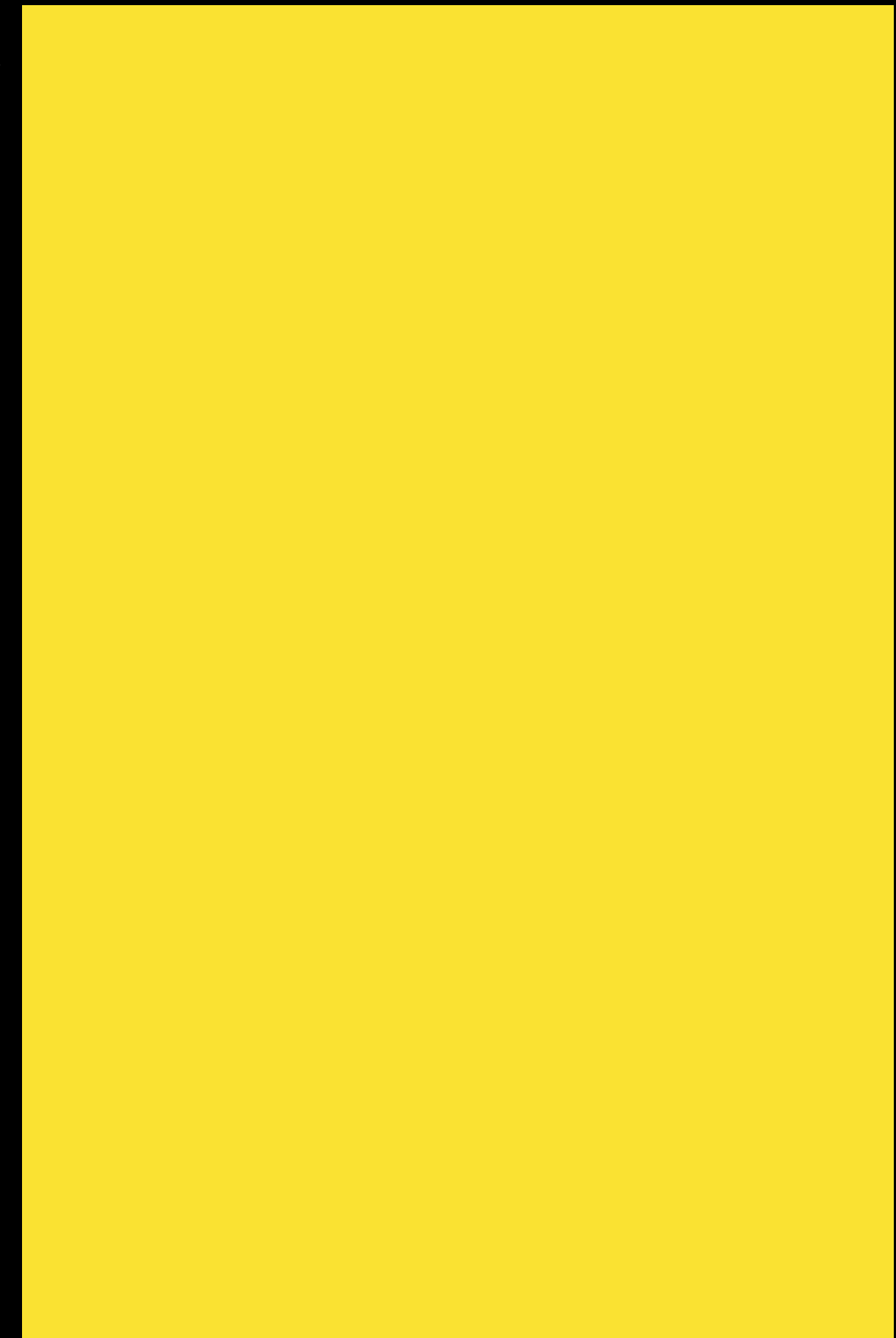


## epilogue

```
mov    rsp, rbp
pop    rbp
ret
```

rsp →

Stack



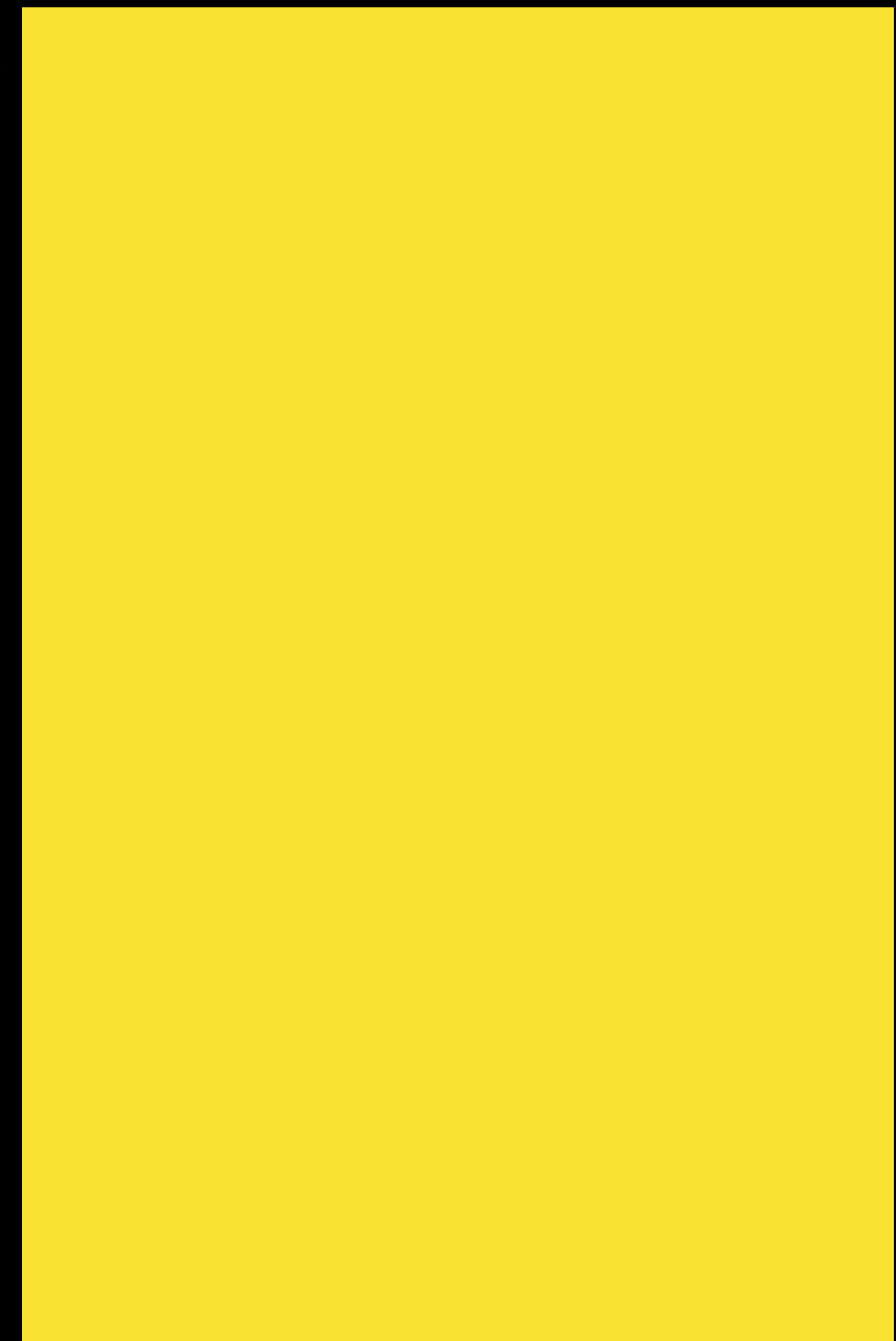
# Memory...

...needs to be function local



Stack

rsp →



**epilogue**

```
mov    rsp, rbp
pop    rbp
ret
```

return to caller

# Memory...

...needs to be function local



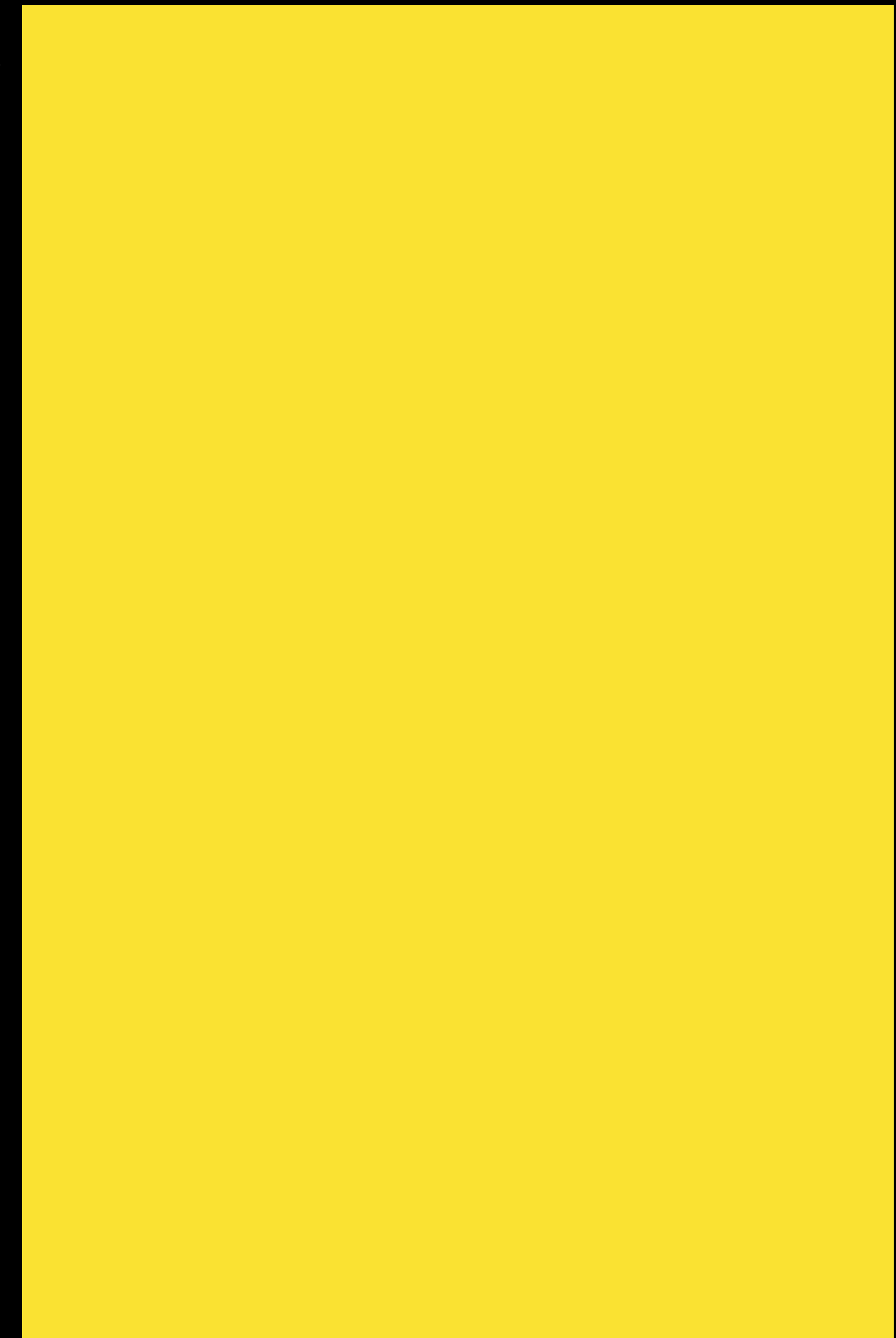
## epilogue

```
mov    rsp, rbp
pop    rbp
ret
```

**After a function call returns,  
the stack of the previous  
function remains unchanged**

rsp →

Stack



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...needs to be function local

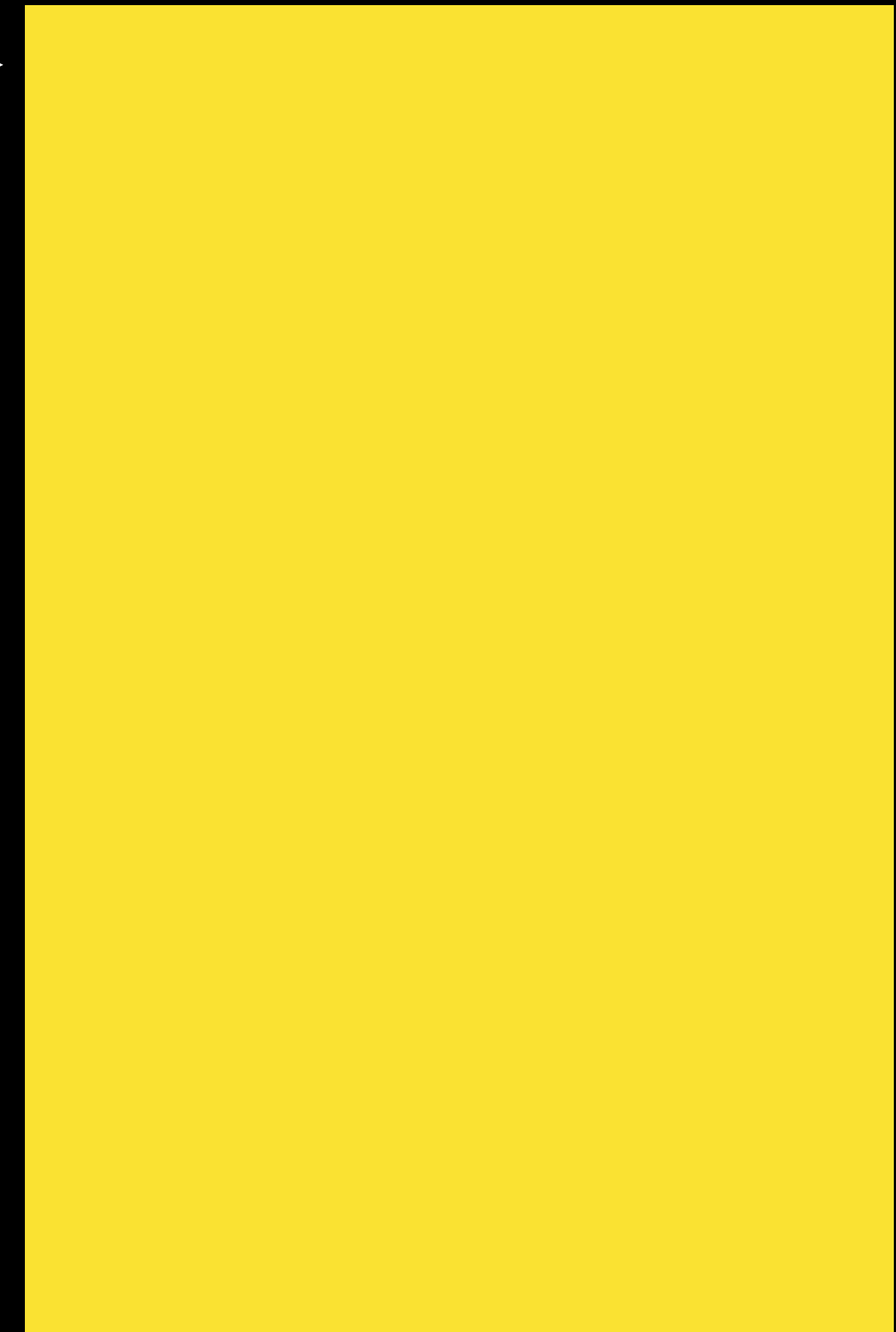


## epilogue

```
mov    rsp, rbp
pop    rbp
ret
```

rsp →

Stack



**After a function call returns,  
the stack of the previous  
function remains unchanged**

**rsp/rbp should remain exactly  
as they were before and after  
the function call**

# Decompilation...

... can fail



- Certain tricks can be used to trip up a decompiler

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- Some decompilers will fail if the previous invariant cannot be statically proven

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# Decompilation...

... can fail



- Certain tricks can be used to trip up a decompiler
- Some decompilers will fail if the previous invariant cannot be statically proven
  - *rsp/rbp should remain exactly as they were before and after the function call*
- Can we prevent this invariant from being proven?

# Decompilation...

... can fail



...prologue

```
        call    MyFunc
        cmp     rax, 0
        jz      CONTINUE
        add     rsp, 4
CONTINUE:
        pop     rax
```

...rest of function

# Decompilation...

## ... can fail



- Let's say that MyFunc always returns 0

...prologue

```
    call    MyFunc
    cmp     rax, 0
    jz      CONTINUE
    add     rsp, 4
CONTINUE:
    pop     rax
```

...rest of function

# Decompilation...

## ... can fail



- Let's say that MyFunc always returns 0
  - rax=0 after call

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    add     rsp, 4
CONTINUE:
    pop     rax
```

...rest of function

# Decompilation...

## ... can fail



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    call    MyFunc
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CONTINUE:
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```

...rest of function

- Let's say that MyFunc always returns 0
  - `rax=0` after call
- Does the decompiler know that MyFunc always returns 0?

# Decompilation...

## ... can fail



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    call    MyFunc
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CONTINUE:
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...rest of function

- Let's say that MyFunc always returns 0
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- Does the decompiler know that MyFunc always returns 0?
  - hint: depending on the decompiler, maybe/maybe not

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## ... can fail



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CONTINUE:
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```

...rest of function

- Let's say that MyFunc always returns 0
  - `rax=0` after call
- Does the decompiler know that MyFunc always returns 0?
  - hint: depending on the decompiler, maybe/maybe not
- Decompiler cannot statically prove that `add rsp, 4` will not get executed

# Decompilation...

... can fail

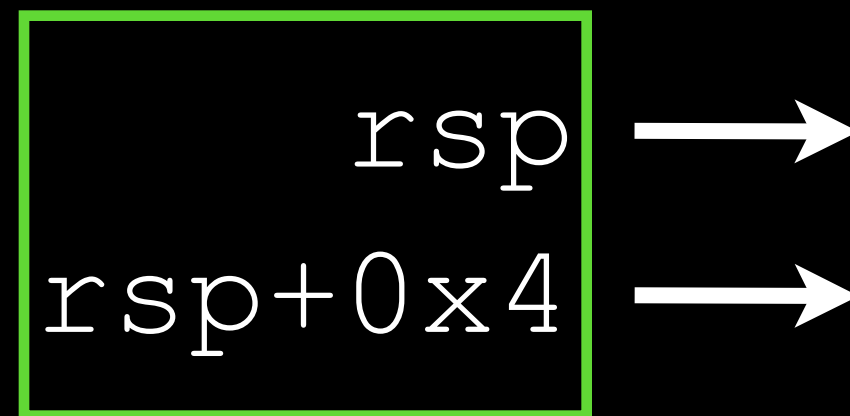


After a function call returns,  
the stack of the previous  
function remains unchanged

rsp/rbp should remain exactly  
as they were before and after  
the function call

Decompiler may not be able to prove  
statically which value will be held in  
the stack pointer after function exit

Stack



# Decompilation...

... can fail in multiple ways



- There are an endless list of anti-decompiler tricks

# Decompilation...

... can fail in multiple ways



- There are an endless list of anti-decompiler tricks
- No time to cover all of them

# Decompilation...

... can fail in multiple ways



- There are an endless list of anti-decompiler tricks
- No time to cover all of them
- Expect some on the challenges :)



Have a break :)

# Debuggers...

... let you have full introspection



- A tool that lets you control/observe a program while it runs

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- Inspect what's happening “under the hood”



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- Step through instructions (or source lines) one at a time

# Debuggers...

... let you have full introspection



- A tool that lets you control/observe a program while it runs
- Inspect what's happening “under the hood”
- Step through instructions (or source lines) one at a time
- Observe memory, registers, local variables, ...

# Debuggers...

... are really useful

- Static reversing only gets you so far



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- Debuggers let you examine runtime behaviour

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  - See what functions return

# Debuggers...

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- Static reversing only gets you so far
- Debuggers let you examine runtime behaviour
  - See real values of variables
  - See what functions return
  - Understand which branches your input causes the program to take



# GDB Demo

# GDB

## Cheat sheet



<b>Start</b>	<code>gdb ./my_prog</code>	<b>Show backtrace/frame</b>	<code>(gdb) bt</code> <code>(gdb) frame</code>
<b>Set breakpoint</b>	<code>(gdb) b my_func</code> <code>(gdb) b *0x400284</code>	<b>Inspect registers</b>	<code>(gdb) info reg</code>
<b>Step into</b>	<code>(gdb) step</code> <code>(gdb) s</code>	<b>Inspect memory</b>	<code>(gdb) x/s <i>address</i></code> <code>(gdb) x/20gx \$rsp</code>
<b>Step over</b>	<code>(gdb) next</code> <code>(gdb) n</code>	<b>Disassemble</b>	<code>(gdb) disas</code> <code>(gdb) x/20i \$rip</code>
<b>Step until function exit</b>	<code>(gdb) finish</code> <code>(gdb) fin</code>	<b>Watch variable/memory</b>	<code>(gdb) watch <i>varname</i></code>
<b>Continue until program stops (i.e., breakpoint)</b>	<code>(gdb) continue</code> <code>(gdb) c</code>	<b>Quit</b>	<code>(gdb) q</code>

# Anti debug tricks

## cat and mouse



- There are multiple ways for a program to detect if it is being run under a debugger

# Anti debug tricks

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- There are multiple ways for a program to detect if it is being run under a debugger
- On Linux, one such technique is using `ptrace`

# Anti debug tricks

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- There are multiple ways for a program to detect if it is being run under a debugger
- On Linux, one such technique is using `ptrace`

```
long ptrace(enum __ptrace_request op, pid_t pid,  
            void *addr, void *data);
```

The **`ptrace()`** system call provides a means by which one process (the "tracer") may observe and control the execution of another process (the "tracee"), and examine and change the tracee's memory and registers. It is primarily used to implement breakpoint debugging and system call tracing.

# Anti debug tricks

## cat and mouse



```
#include <iostream>
#include <sys/ptrace.h>

int main()
{
    if (ptrace(PTRACE_TRACEME, 0, 1, 0) == -1)
    {
        std::cout << "\033[1;31mgo away\033[0m" << std::endl;
        return 1;
    }

    std::cout << "very very secret stuff do not debug thank you" << std::endl;
    return 0;
}
```

16,0-1 All

# Anti debug tricks

## cat and mouse



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~
~
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16,0-1 All

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16,0-1

All

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```
$ c++ ptrace_demo.cpp -o ptrace_demo
$ ./ptrace_demo
very very secret stuff do not debug thank you
$ gdb ./ptrace_demo
Reading symbols from ./ptrace_demo...
(No debugging symbols found in ./ptrace_demo)
(gdb) r
Starting program: /tmp/ptrace_demo
Function(s) ^std::(move|forward|as_const|(__)?addressof) will be skipped when stepping
.
Function(s) ^std::(shared|unique)_ptr<.*>::(get|operator) will be skipped when stepping
g.
Function(s) ^std::(basic_string|vector|array|deque|(forward_)?list|(unordered_|flat_)?
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ed when stepping.
[Thread debugging using libthread_db enabled]
Using host libthread_db library "/lib/x86_64-linux-gnu/libthread_db.so.1".
go away
[Inferior 1 (process 2320934) exited with code 01]
(gdb)
```

# Anti debug tricks

## cat and mouse



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16,0-1

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## cat and mouse

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# Anti debug tricks

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  - Runtime `.text` CRC computation

# Anti debug tricks

## cat and mouse



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  - `int3` (debugger trap)

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- Many such techniques
  - Runtime `.text` CRC computation
  - `int3` (debugger trap)
  - `/proc/self/status` (TracerPid)
- This talk is already very long, so I'll leave learning these as an exercise to you

# Anti debug tricks

## cat and mouse



- Many such techniques
  - Runtime `.text` CRC computation
  - `int3` (debugger trap)
  - `/proc/self/status` (TracerPid)
- This talk is already very long, so I'll leave learning these as an exercise to you
- You can always find a way around these (i.e., patch the program)

# Packers...

... pack stuff

- Executables can get pretty big



# Packers...

... pack stuff



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- **Packers** exist to compress an executable into a smaller binary

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- Resulting binary is compressed payload + decompressor
- At runtime, decompresses payload and jumps to it
- Common in malware and some “commercial software protection solutions”
- You can typically recognise the use of a packer from `strings`
  - “This file is packed with the UPX executable packer”

# Packers...

... pack stuff



Original Program

# Packers...

... pack stuff

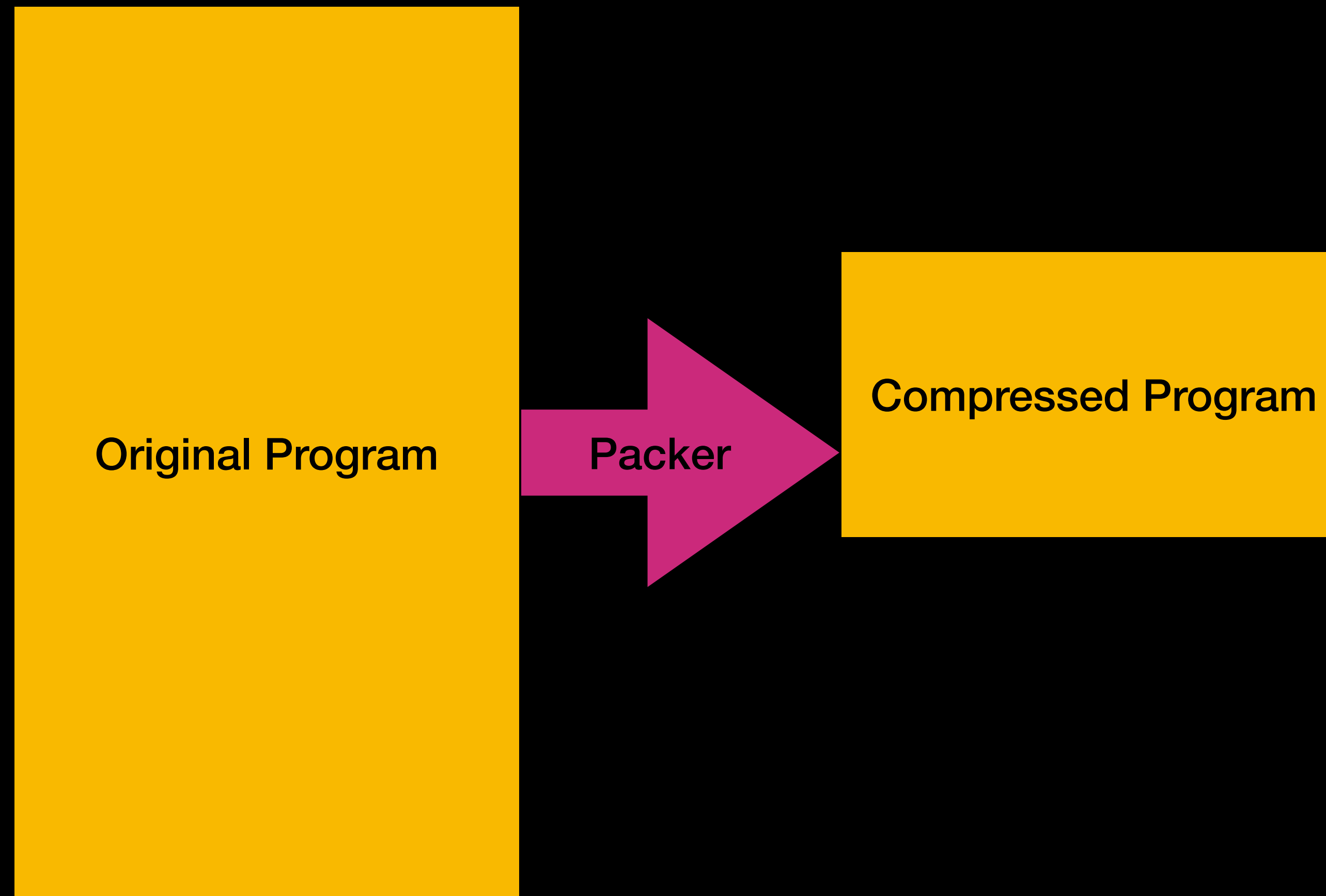


Original Program

Packer

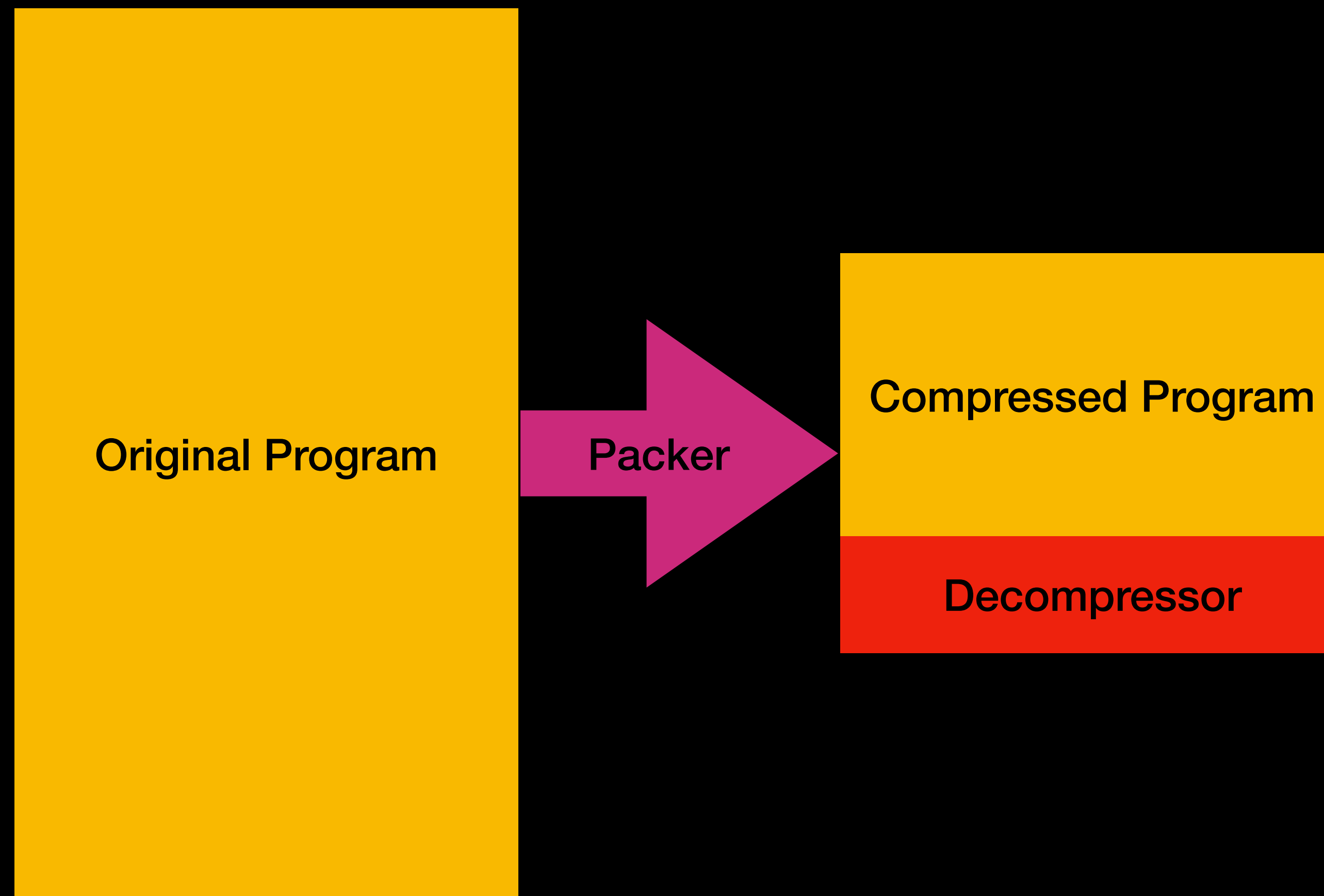
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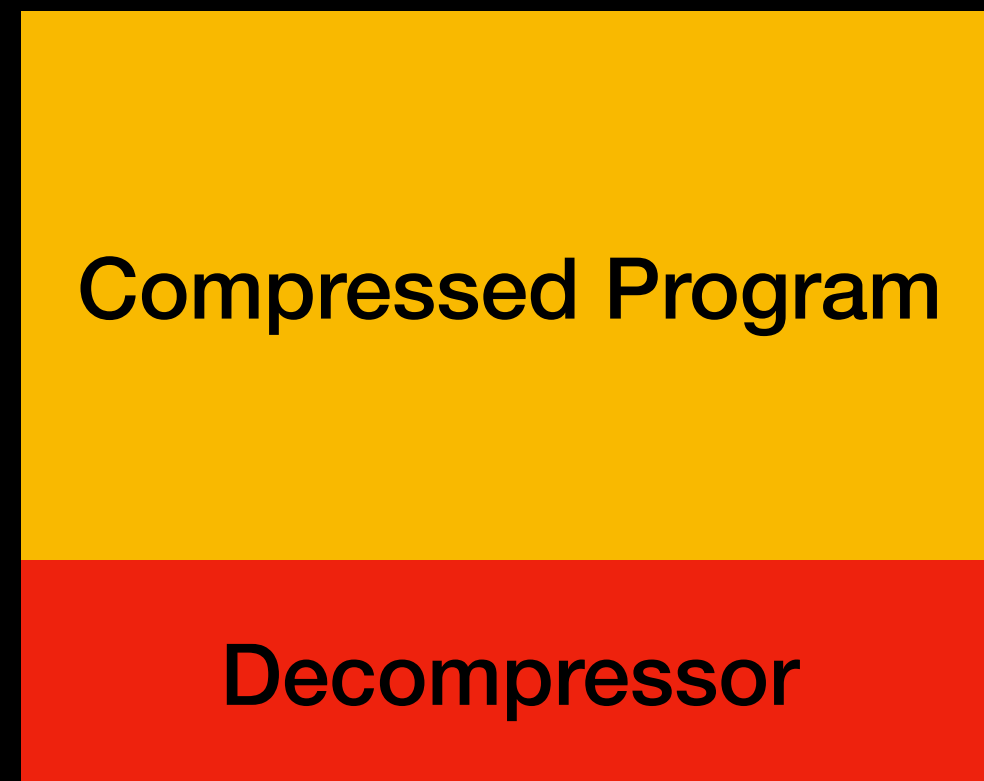
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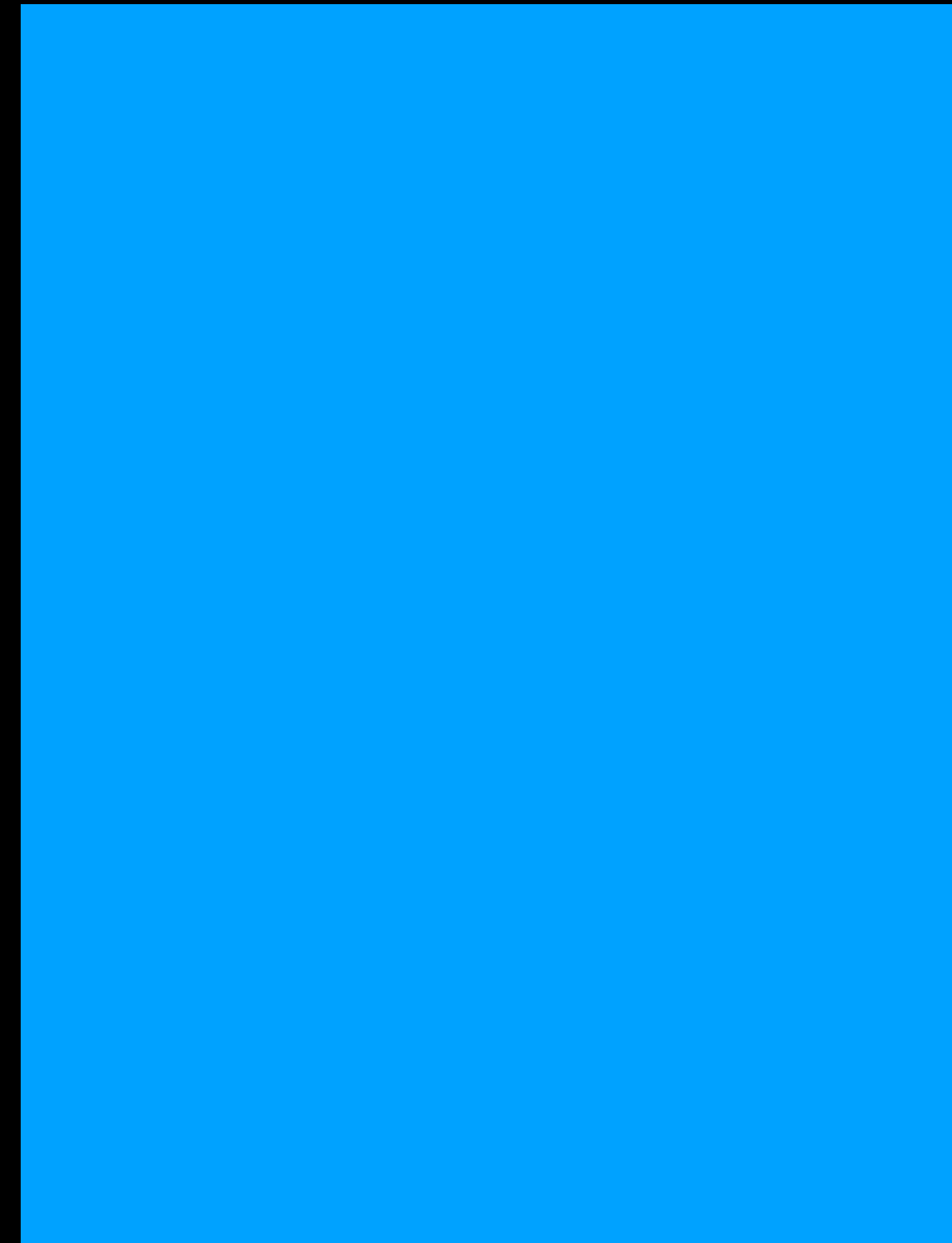


# Packers...

... also unpack stuff

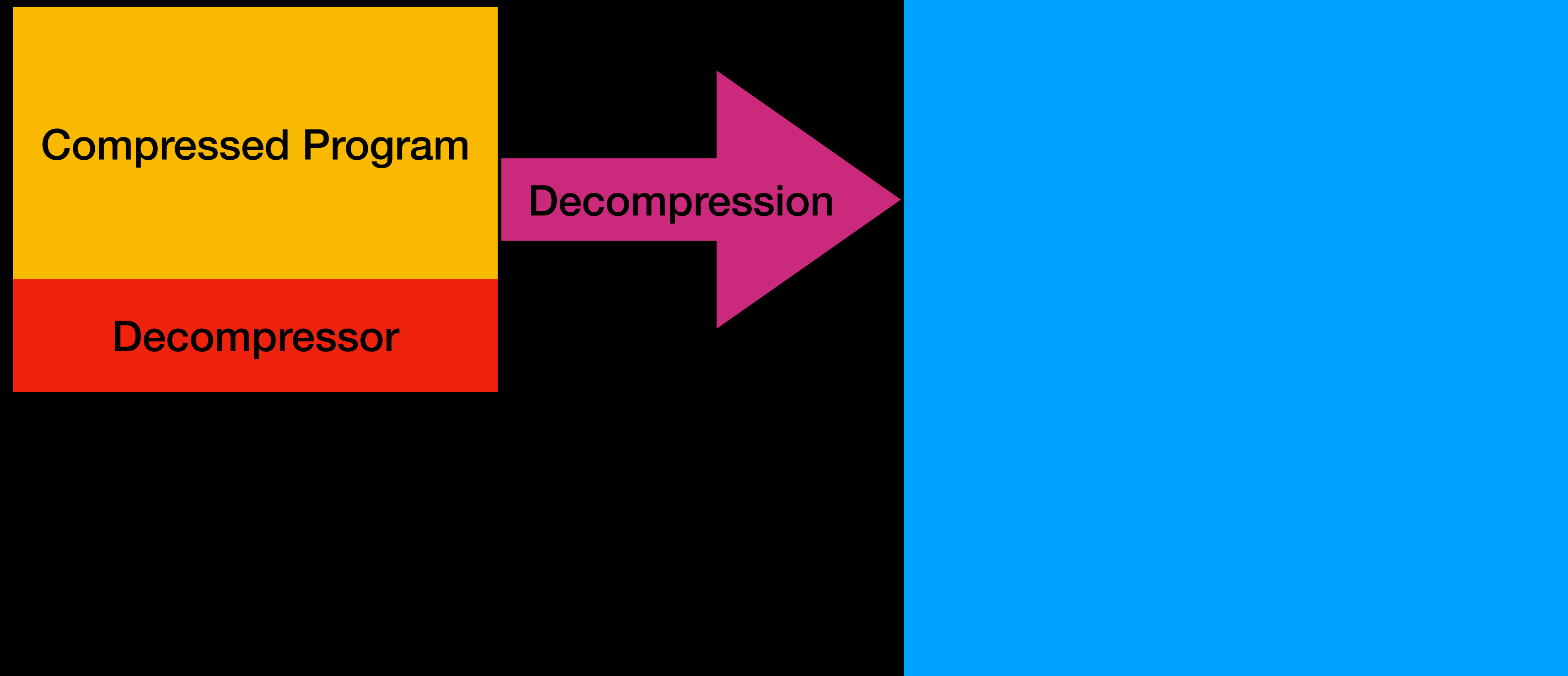


Memory / RAM



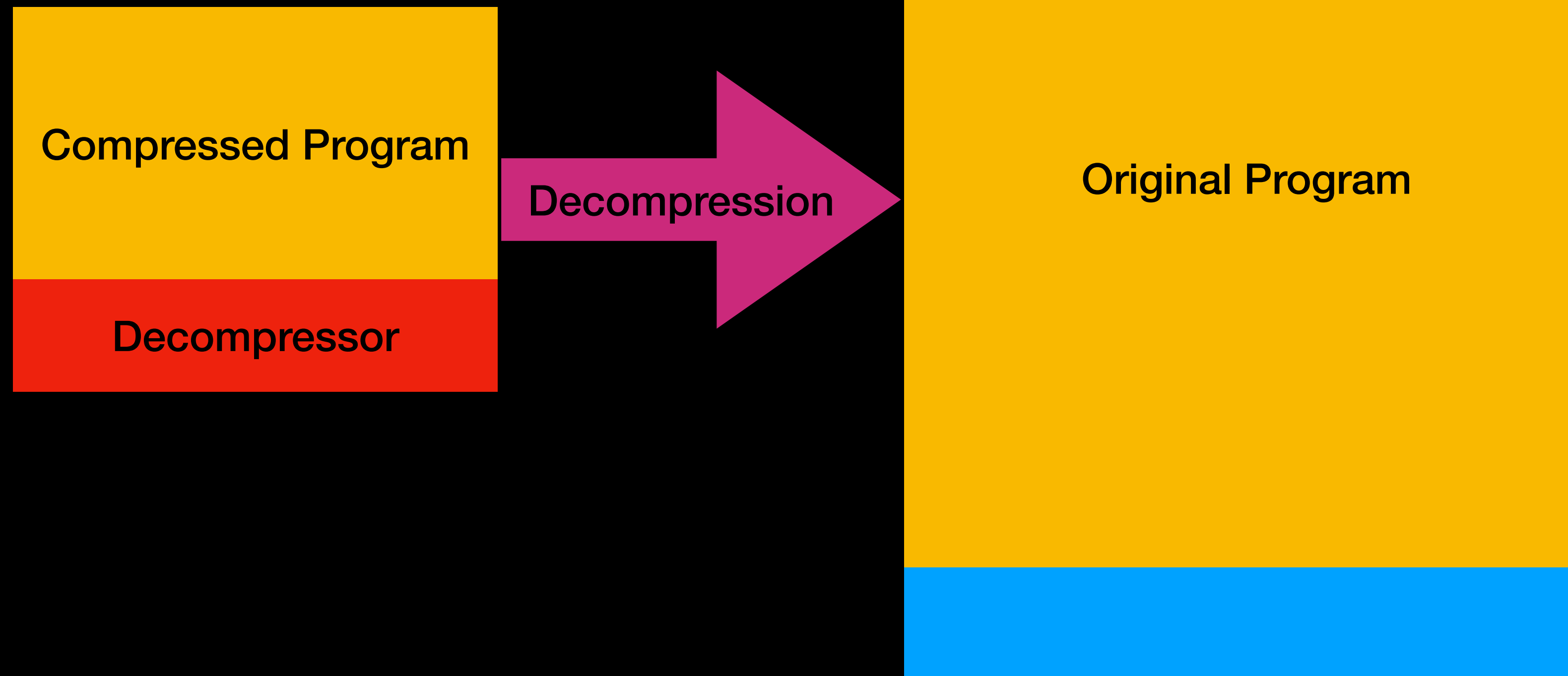
# Packers...

... also unpack stuff



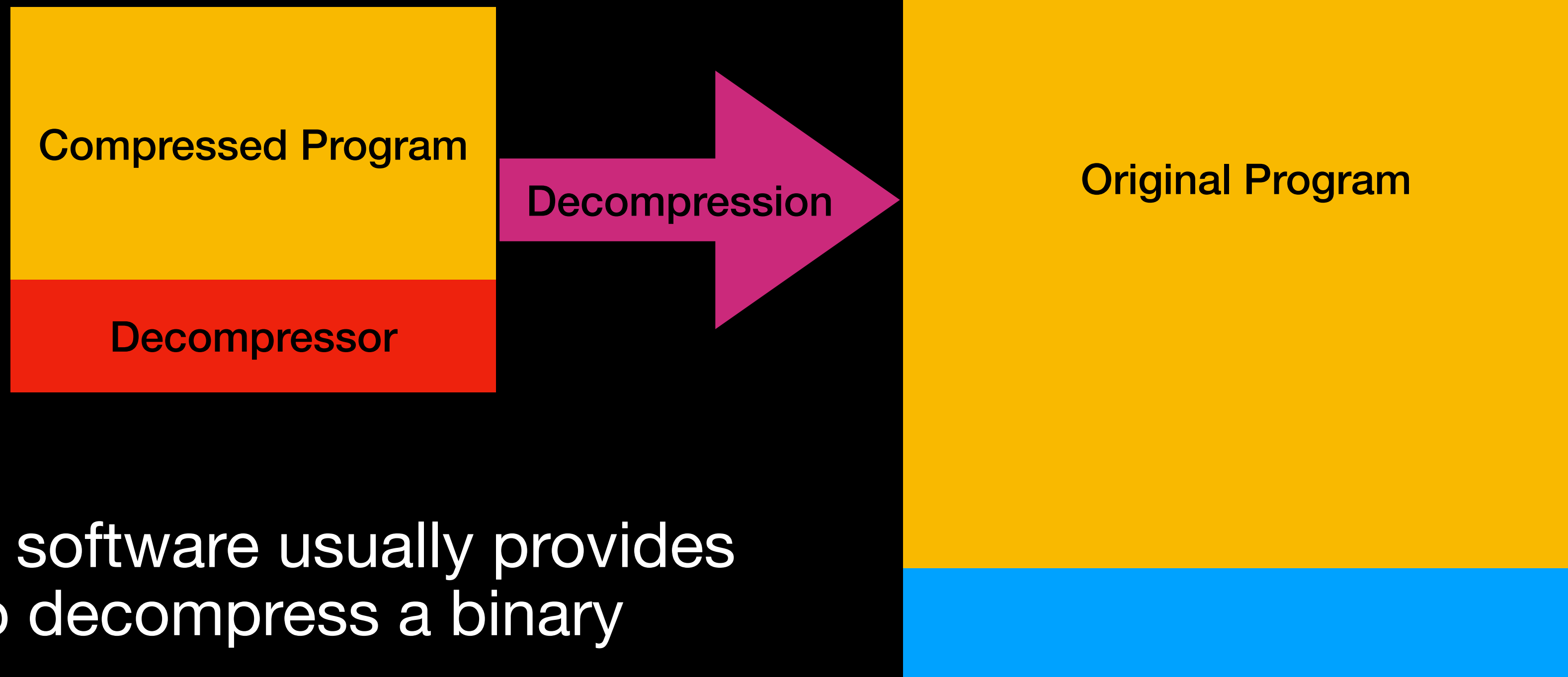
# Packers...

... also unpack stuff



# Packers...

... also unpack stuff



Packing software usually provides  
a way to decompress a binary

i.e. `upx -d`

# Packers...

... pack and unpack things

- A lot of packers exist



# Packers...

... pack and unpack things

- A lot of packers exist
  - Yoda



# Packers...

... pack and unpack things

- A lot of packers exist
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  - UPX



# Packers...

## ... pack and unpack things



- A lot of packers exist
  - Yoda
  - UPX
  - 20to4

# Packers...

## ... pack and unpack things



- A lot of packers exist
  - Yoda
  - UPX
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  - eXPressor

# Packers...

... pack and unpack things



- A lot of packers exist
  - Yoda
  - UPX
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  - eXPressor
- Some may offer password protection, encryption, etc

# Packers...

... pack and unpack things



- A lot of packers exist
  - Yoda
  - UPX
  - 20to4
  - eXPressor
- Some may offer password protection, encryption, etc
- Another “exercise to the reader” to go and learn more

# Obfuscaters...

... make your life harder



- Obfuscation is the process of transforming code so it's harder to read but still does the same thing

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# Obfuscaters...

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- Goal: frustrate reverse engineers (you!)
- Commonly used in malware, DRM, and CTFs :^)

# Obfuscaters...

... why?

- Hide algorithms, secret constants (decryption keys, flags)



# Obfuscators...

... why?

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- Slow down reverse engineers



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- Avoid threat detection (less common)



# Obfuscators...

... why?



- Hide algorithms, secret constants (decryption keys, flags)
- Slow down reverse engineers
- Avoid threat detection (less common)
- Increase cost of defence

# Obfuscaters...

... how? - control flow flattening

- Turn neat control flow structures into spaghetti mess



# Obfuscators...

... how? - control flow flattening

- Turn neat control flow structures into spaghetti mess

```
if (x == 5) return 1;  
else return 0;
```



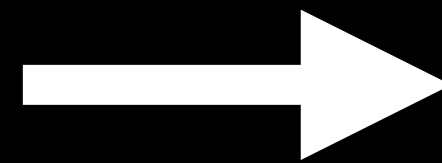
# Obfuscators...

... how? - control flow flattening



- Turn neat control flow structures into spaghetti mess

```
if (x == 5) return 1;  
else return 0;
```



```
state = 0;  
while (1) {  
    switch (state) {  
        case 0:  
            if (x == 5) state = 1;  
            else state = 2;  
            break;  
        case 1: return 1;  
        case 2: return 0;  
    }  
}
```

# Obfuscators...

... how? - control flow flattening



- Turn neat control flow structures into spaghetti mess

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Both of these return 1 if  $x == 5$

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But this is much harder  
to reason about

# Obfuscaters...

... how? - string encryption

- Encrypt all strings, and only decrypt when that string is used



# Obfuscaters...

... how? - string encryption



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#include <stdio.h>
int main() {
    printf("The secret is RISC{obf_string}\n");
    return 0;
}
```

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#include <stdio.h>
#include <string.h>

void deobf(char *s, int key) {
    for (int i = 0; s[i]; i++) {
        s[i] ^= key;
    }
}

int main() {
    char flag[] = "\x07\x1c\x06\x16.:73\n&!'<;2(";
    deobf(flag, 0x55);
    printf("The secret is %s\n", flag);
    return 0;
}
```

# Obfuscators...

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... how?

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- Can range from simple string encryption...



# Obfuscaters...

... how?



- Non exhaustive list, of course
- Can range from simple string encryption...
- ... to reducing the entire program down to a single instruction (MOV)

# Obfuscaters...

## ... how?



- No
- C
- ...

```
<is_prime>:
push    ebp
mov     ebp,esp
sub     esp,0x10
cmp     DWORD PTR [ebp+0x8],0x1
jne     8048490 <is_prime+0x13>
mov     eax,0x0
jmp     80484cf <is_prime+0x52>
cmp     DWORD PTR [ebp+0x8],0x2
jne     804849d <is_prime+0x20>
mov     eax,0x1
jmp     80484cf <is_prime+0x52>
mov     DWORD PTR [ebp-0x4],0x2
jmp     80484be <is_prime+0x41>
mov     eax,DWORD PTR [ebp+0x8]
cdq
idiv    DWORD PTR [ebp-0x4]
mov     eax,edx
test    eax,eax
jne     80484ba <is_prime+0x3d>
mov     eax,0x0
jmp     80484cf <is_prime+0x52>
add     DWORD PTR [ebp-0x4],0x1
mov     eax,DWORD PTR [ebp-0x4]
imul    eax,DWORD PTR [ebp-0x4]
cmp     eax,DWORD PTR [ebp+0x8]
jle     80484a6 <is_prime+0x29>
mov     eax,0x1
leave
ret
```

encryption...

am down to a single instruction (MOV)

# Obfuscaters...

## ... how?



- No
- C
- ...

```
<is_prime>:
push    ebp
mov     ebp,esp
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mov     eax,0x0
jmp     80484cf <is_prime+0x52>
cmp     DWORD PTR [ebp+0x8],0x2
jne     804849d <is_prime+0x20>
mov     eax,0x1
jmp     80484cf <is_prime+0x52>
mov     DWORD PTR [ebp-0x4],0x2
jmp     80484be <is_prime+0x41>
mov     eax,DWORD PTR [ebp+0x8]
cdq
idiv    DWORD PTR [ebp-0x4]
mov     eax,edx
test    eax,eax
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mov     eax,0x0
jmp     80484cf <is_prime+0x52>
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mov     eax,DWORD PTR [ebp-0x4]
imul    eax,DWORD PTR [ebp-0x4]
cmp     eax,DWORD PTR [ebp+0x8]
jle     80484a6 <is_prime+0x29>
mov     eax,0x1
leave
ret
```

encrypti  
am dov

```
mov     dl,BYTE PTR ds:0x81fc4d0
mov     eax,DWORD PTR [eax*4+0x81fbc30]
mov     eax,DWORD PTR [eax+edx*4+0x81fac90]
mov     ds:0x81fc55f,al
mov     BYTE PTR ds:0x81fc4d0,ah
mov     DWORD PTR ds:0x81fc4d0,0x0
mov     eax,ds:0x81fc55c
mov     ds:0x81fc4c0,eax
mov     eax,ds:0x81fc554
mov     ds:0x81fc4c4,eax
mov     eax,0x0
mov     ecx,0x0
mov     DWORD PTR ds:0x81fc4d0,0x1
mov     ax,ds:0x81fc4c0
mov     cx,WORD PTR ds:0x81fc4c4
mov     cx,WORD PTR [ecx*2+0x8167520]
mov     edx,DWORD PTR [eax*4+0x8067400]
mov     edx,DWORD PTR [edx+ecx*4]
mov     edx,DWORD PTR [edx*4+0x8067400]
mov     ecx,DWORD PTR ds:0x81fc4d0
mov     edx,DWORD PTR [edx+ecx*4]
mov     WORD PTR ds:0x81fc560,dx
mov     ecx,DWORD PTR ds:0x81fc4d0
mov     edx,DWORD PTR [edx+ecx*4]
mov     edx,DWORD PTR [edx*4+0x8067400]
mov     ecx,DWORD PTR ds:0x81fc4d0
mov     edx,DWORD PTR [edx+ecx*4]
mov     WORD PTR ds:0x81fc562,dx
mov     DWORD PTR ds:0x81fc4ce,edx
mov     eax,0x0
mov     al,ds:0x81fc4d0
mov     al,BYTE PTR [eax+0x80535d0]
mov     ds:0x81fc560,eax
mov     eax,ds:0x81fc560
mov     edx,DWORD PTR [eax*4+0x81fc504]
mov     DWORD PTR ds:0x81fc5a4,edx
mov     edx,DWORD PTR [eax*4+0x81fc504]
mov     DWORD PTR ds:0x81fc5ac,edx
mov     eax,ds:0x81fc5a4
mov     eax,DWORD PTR [eax]
mov     ds:0x81fc500,eax
mov     eax,ds:0x81fc500
mov     ds:0x81fc4c0,eax
mov     eax,ds:0x81fc554
mov     ds:0x81fc4c4,eax
mov     eax,0x0
mov     ecx,0x0
mov     DWORD PTR ds:0x81fc4d0,0x1
mov     ax,ds:0x81fc4c0
mov     cx,WORD PTR ds:0x81fc4c4
mov     cx,WORD PTR [ecx*2+0x8167520]
mov     edx,DWORD PTR [eax*4+0x8067400]
mov     edx,DWORD PTR [edx+ecx*4]
mov     edx,DWORD PTR [edx*4+0x8067400]
mov     ecx,DWORD PTR ds:0x81fc4d0
mov     edx,DWORD PTR [edx+ecx*4]
mov     WORD PTR ds:0x81fc560,dx
mov     DWORD PTR ds:0x81fc4ce,edx
mov     ax,ds:0x81fc4c2
mov     cx,WORD PTR ds:0x81fc4c6
```

```
mov     ds:0x81fc500,eax
mov     eax,ds:0x81fc500
mov     ds:0x81fc4c0,eax
mov     eax,ds:0x81fc554
mov     ds:0x81fc4c4,eax
mov     eax,0x0
mov     ecx,0x0
mov     DWORD PTR ds:0x81fc4d0,0x1
mov     ax,ds:0x81fc4c0
mov     cx,WORD PTR ds:0x81fc4c4
mov     cx,WORD PTR [ecx*2+0x8167520]
mov     edx,DWORD PTR [eax*4+0x8067400]
mov     edx,DWORD PTR [edx+ecx*4]
mov     edx,DWORD PTR [edx*4+0x8067400]
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mov     edx,DWORD PTR [edx+ecx*4]
mov     WORD PTR ds:0x81fc500,dx
mov     DWORD PTR ds:0x81fc4ce,edx
mov     ax,ds:0x81fc4c2
mov     cx,WORD PTR ds:0x81fc4c6
mov     cx,WORD PTR [ecx*2+0x8167520]
mov     edx,DWORD PTR [eax*4+0x8067400]
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mov     edx,DWORD PTR [edx*4+0x8067400]
mov     ecx,DWORD PTR ds:0x81fc4d0
mov     edx,DWORD PTR [edx+ecx*4]
mov     WORD PTR ds:0x81fc502,dx
mov     DWORD PTR ds:0x81fc4ce,edx
mov     eax,ds:0x81fc5a4
mov     edx,DWORD PTR ds:0x81fc500
mov     DWORD PTR [eax],edx
mov     eax,ds:0x81fc5ac
mov     eax,DWORD PTR [eax]
mov     ds:0x81fc57c,eax
mov     eax,0x0
mov     al,ds:0x81fc57e
mov     al,BYTE PTR [eax+0x8055a94]
mov     ds:0x81fc57e,al
mov     eax,ds:0x81fc5ac
mov     edx,DWORD PTR ds:0x81fc57c
mov     DWORD PTR [eax],edx
mov     edx,0x0
mov     dl,BYTE PTR ds:0x81fc551
mov     eax,DWORD PTR [edx*4+0x8055660]
mov     ds:0x81fc4d0,eax
mov     eax,0x0
mov     edx,0x0
mov     al,ds:0x81fc55c
mov     dl,BYTE PTR ds:0x81fc4d0
mov     eax,DWORD PTR [eax*4+0x81fbc30]
mov     eax,DWORD PTR [eax+edx*4+0x81fac90]
mov     ds:0x81fc55c,al
mov     BYTE PTR ds:0x81fc4d0,ah
mov     eax,0x0
mov     edx,0x0
mov     al,ds:0x81fc55d
mov     dl,BYTE PTR ds:0x81fc4d0
mov     eax,DWORD PTR [eax*4+0x81fbc30]
mov     eax,DWORD PTR [eax+edx*4+0x81fac90]
mov     ds:0x81fc55d,al
mov     BYTE PTR ds:0x81fc4d0,ah
mov     eax,0x0
mov     edx,0x0
mov     al,ds:0x81fc55e
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```
mov     ds:0x81fc4d0,eax
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mov     al,ds:0x81fc55c
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mov     edx,0x0
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mov     eax,DWORD PTR [eax+edx*4+0x81fac90]
mov     ds:0x81fc55e,al
mov     BYTE PTR ds:0x81fc4d0,ah
mov     eax,0x0
mov     edx,0x0
mov     al,ds:0x81fc55f
mov     dl,BYTE PTR ds:0x81fc4d0
mov     eax,DWORD PTR [eax*4+0x81fbc30]
mov     eax,DWORD PTR [eax+edx*4+0x81fac90]
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mov     ecx,0x0
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```



## M/o/Vfuscator

- No
- Ca
- ...

# Obfuscaters...

... how?

- Non exhaustive list, of course



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# Bonus: SMT solvers



- Many reversing problems boil down to “Find input  $X$  that satisfies constraints  $Y$ ”

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- Simple example:

```
If (a * 3 + b == 42 && (a ^ b) == 7) {  
    win();  
}
```

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- SMT solvers find solutions to these equations
  - i.e., finding X such that Y

Exercise for the reader :)

- Simple example:

```
If (a * 3 + b == 42 && (a ^ b) == 7) {  
    win();  
}
```



Questions?

# Tools that you should go and learn



- `strings`
- `Binary Ninja`
- `objdump`
- `binwalk`
- `file`
- `gdb`
- `nc`



# Feedback



<https://forms.office.com/r/3L9rMBd2iQ>



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